

Center for Quantum Networks The development of this short course was supported primarily by the Engineering Research Center (ERC) Program of the National Science Foundation under NSF Cooperative Agreement No. 1941583. Any opinions, findings and conclusions or recommendations expressed in our course material are those of the instructors and do not necessarily reflect those of the National Science Foundation.

Information in a Photon

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CQN Winter School on Quantum Networks January 5, 2023

Funded by National Science Foundation Grant #1941583















Quantum information processing



- Communications
 - Deep-space lasercom
 - Quantum security: QKD, covert communications
- Computing
 - Factoring [breaks RSA, Diffie-Hellman]
 - Search and optimization
 - Blind quantum computing
 - Multiparty privacy-preserving computations
 - Simulations of complex quantum systems
- Sensing
 - Passive imaging (astronomy, microscopy, SDA)
 - Active photonic sensors (RF photonic antennas, fiber-optic gyroscopes beam deflectometers, surface topography, endoscopy, lidars, ...)
 - Gravitational wave sensors (LIGO squeezed light)
 - Magnetic field sensing (Neuronal imaging, chip testing)



Photonic quantum information processing

 Wherever light gets used in encoding, extracting, carrying or processing information
[communications, sensing, imaging, computing, simulations, ...], what is the best performance permissible by the laws of quantum physics? Quantum optics meets Information theory

Quantum optics – quantum theory of light and its interaction with matter

- Non-linear optics
- Atomic systems
- Many-body physics

Information theory – quantifying "information" in the context of communications, sensing and computing

- Detection theory
- Estimation theory
- Data compression
- Channel coding



Course objective

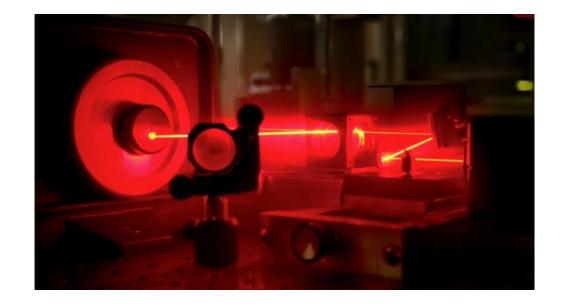
- Introduction to some of the mathematical tools necessary to understand the quantum representation of classical laser light, and certain *non-classical* states of light, including single-photon and multi-photon states and *squeezed* states of light
- Using these tools to understand how to manipulate (classical and quantum) light using *interference*
- Using basic tools from quantum information theory to uncover applications of quantum techniques (sources of light, ways to manipulate light, and detection schemes) for improved performance in encoding, transmission and processing of information



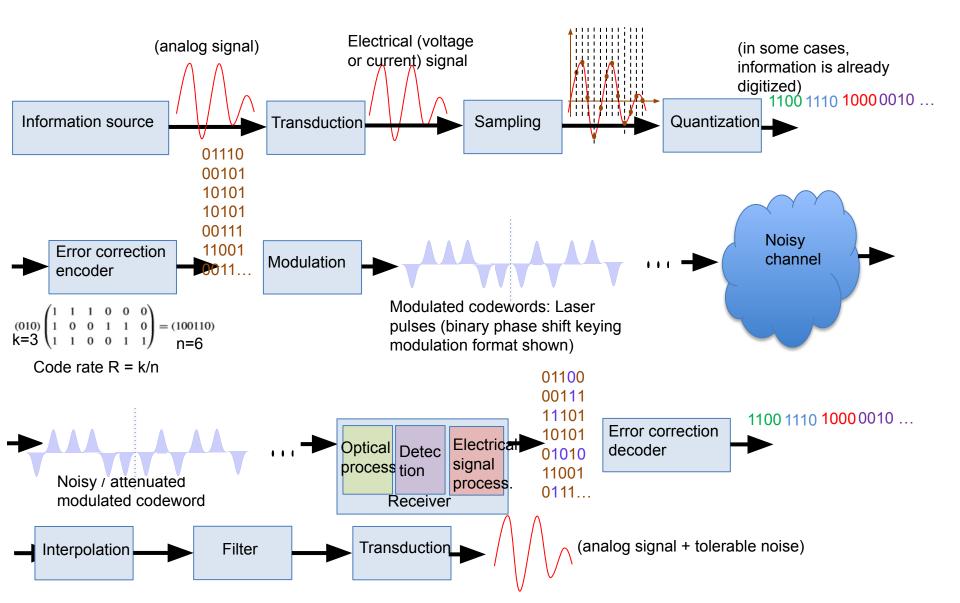
- Module 1: Quantum limits of information encoded in laser light [Saikat Guha]
- **Module 2:** Quantum information advantage arising from interfering photons [Christos Gagatsos]
- Multiple choice questions interspersed through the course, to be answered through zoom polls
- Post-lecture survey



Module 1: Quantum limits of information encoded in laser light



Digital (classical) communications

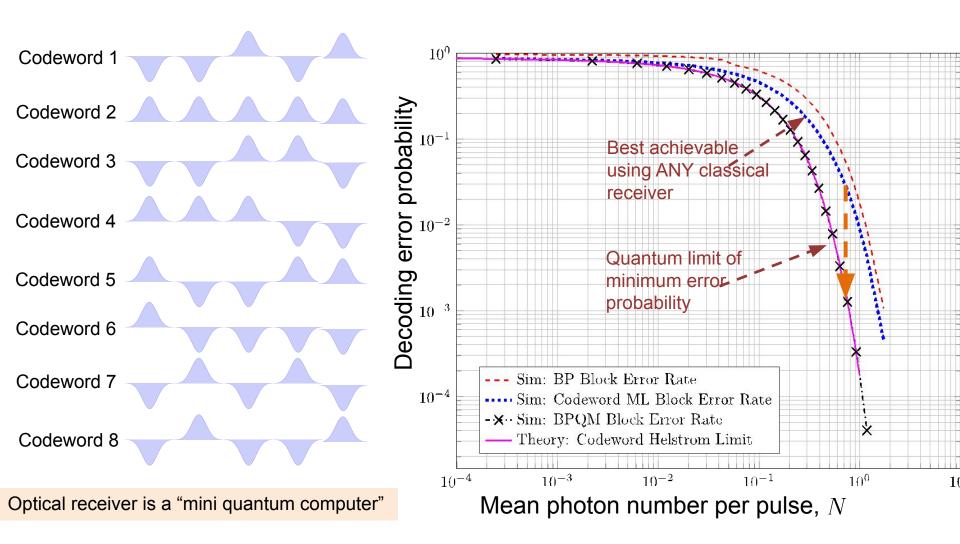


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"Quantum limit" of an optical receiver

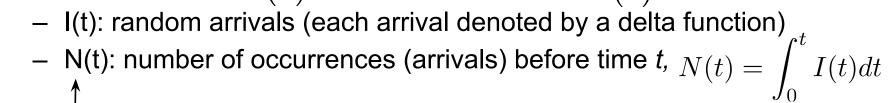


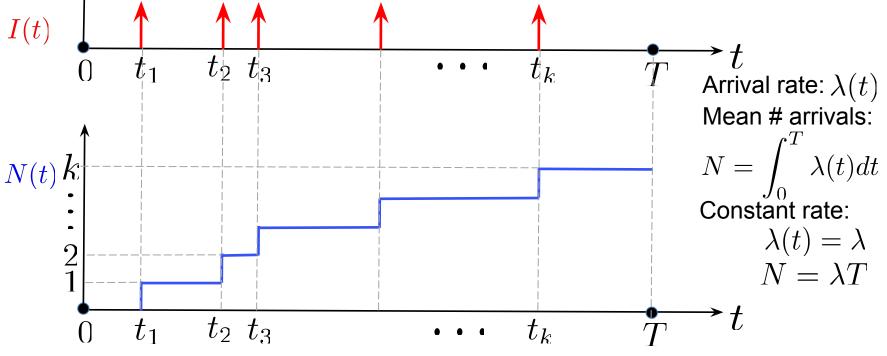
Silva, SG, Dutton, Phys. Rev. A 87, 052320 (2013) Rengaswamy, Seshadreesan, SG, Pfister, Nature npj Quantum Inf 7, 97 (2021) Delaney, Seshadreesan, MacCormack, Galda, SG, Narang, Phys. Rev. A 106, 032613 (2022)

Random process describing *Discrete* events in *continuous* time and/or space



- Where such random processes might occur
 - Spiking patterns in neurons, bank teller queue, photon detection, ...
- Arrival process, I(t); counting process, N(t)

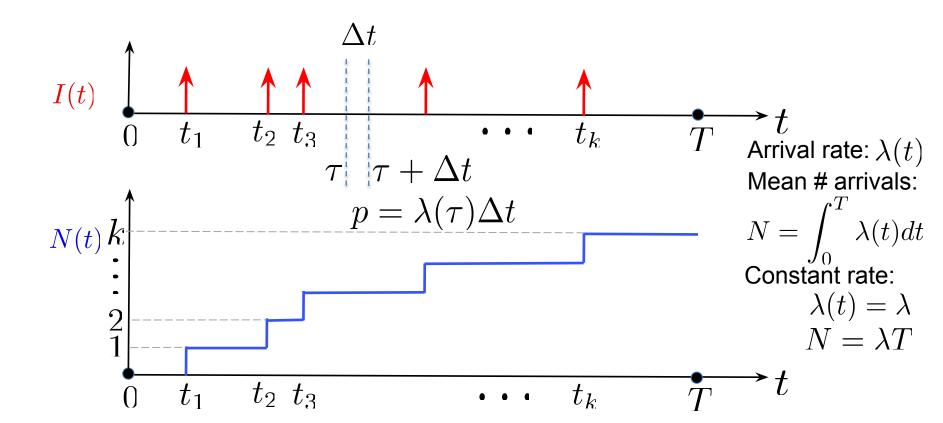




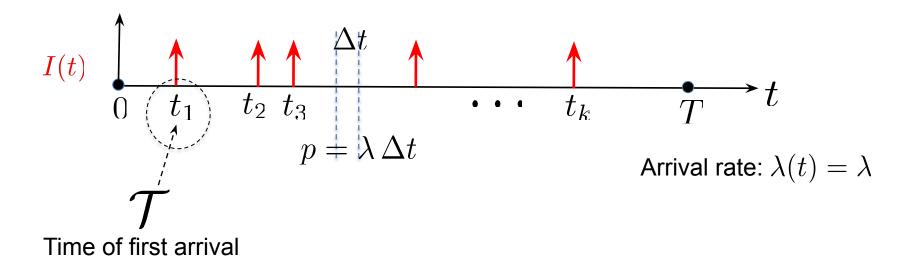


1. Probability of an arrival in a tiny time step is equal to the rate of the arrival process (at that time) times the size of the time increment

2. Probabilities of arrivals in disjoint time steps are statistically independent



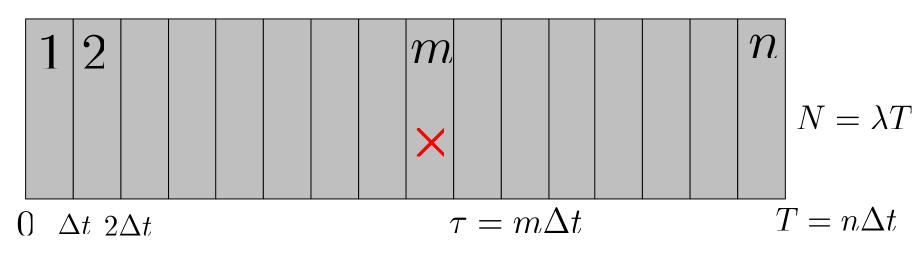
Probability distribution of inter-arrival time



Probability distribution of the time of first arrival, $P_{\mathcal{T}}(\tau), \, \tau \geq 0$

Probability distribution of the total number of arrivals, $P_K[k], \, k=0,1,\ldots$

Inter-arrival times and number of arrivals



Probability of one arrival in a $\, \Delta t \,$ interval, $p = \lambda \, \Delta t \,$

Let us denote by t, the time of first arrival; $P[t = \tau] = (1 - p)^{m-1}p; \ m = 1, 2, ...$ c.d.f., $F_{\mathcal{T}}(\tau) = P[t \le \tau] = \sum_{j=1}^{m} (1 - p)^{j-1}p = 1 - (1 - p)^m$ $= 1 - (1 - \lambda\Delta t)^{\tau/\Delta t} \xrightarrow{\Delta t \to 0} 1 - e^{-\lambda\tau}$ p.d.f., $P_{\mathcal{T}}(\tau) = \frac{d}{d\tau}F_{\mathcal{T}}(\tau) = \lambda e^{-\lambda\tau}; \ \tau \ge 0$ $\stackrel{m \to \infty}{\longrightarrow} \frac{T}{\tau} = \frac{n}{m}$ Probability distribution of the total number of arrivals K, $P_K[k] = \binom{n}{k} p^k (1 - p)^{n-k}$



Distribution of the number of arrivals

$$\tau = m\Delta t$$
 $T = n\Delta t$

$$p = \lambda \, \Delta t \qquad \qquad N = \lambda T$$

Probability distribution of the total number of arrivals K, $P_K[k] = \binom{n}{k} p^k (1-p)^{n-k}$ $\overrightarrow{\Delta t \to 0}$?

Problem 1: What is the distribution of the total number of arrivals K? (k=0, 1, 2, ...)

A:
$$P_K[k] = (1 - 1/N)^k (1/N)$$

B: $P_K[k] = e^{-N} N^k / k!$
C: $P_K[k] = N^{k^2} - N^{(k+1)^2}$
D: I do not know.

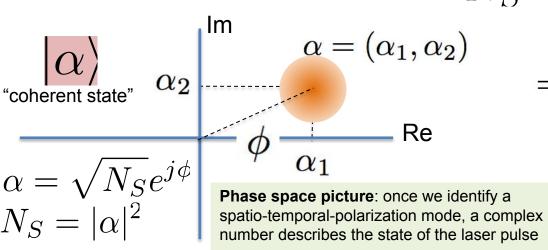


A pulse of laser light

 Quasi-mono-chromatic laser light pulse: in √ (photons/m²-sec) units

$$\tilde{E}(\boldsymbol{r},t) = E(\boldsymbol{r},t)e^{-j\omega_0 t + \phi}, t \in (0,T], \boldsymbol{r} \in \mathcal{A}$$
Spatial and temporal $= \psi(\boldsymbol{r})s(t)e^{-j\omega_0 t + \phi}, \boldsymbol{r} \equiv (x,y)$
dependence may not be
factorable in general Spatial mode Temporal mode

• Mean photon number, $N_S =$



$$\int_{0}^{T} \int_{\mathcal{A}} |\tilde{E}(\boldsymbol{r},t)|^{2} d\boldsymbol{r} dt$$
$$\int_{0}^{T} \int_{\mathcal{A}} |E(\boldsymbol{r},t)|^{2} d\boldsymbol{r} dt$$

No detector can accurately measure the field $E(\pmb{r},t)$



- An optical mode is the "shape" of a confined EM field in space, time and polarization (the three independent degrees of freedom of the photon)
- Time & Frequency are the same degree of freedom (related by Fourier transform)

$$\phi_{\nu}(\boldsymbol{r},t);\,\boldsymbol{r}\in\mathcal{A},t\in[0,T),\nu=1,2$$

$$\int_{\mathcal{A}} \int_{0}^{T} \phi_{\nu}(\boldsymbol{r},t) \phi_{\nu}^{*}(\boldsymbol{r},t) d\boldsymbol{r} dt$$

 $= \int_{\mathcal{A}} \int_{0}^{\mathbf{r}} |\phi_{\nu}(\mathbf{r},t)|^{2} d\mathbf{r} dt = 1$

We will take a mode to be normalized



Orthogonal modes

- Two modes $\phi_{
u}({m r},t)$ and $\psi_{\mu}({m r},t)$ are orthogonal if,

$$\int_{\mathcal{A}} \int_{0}^{T} \phi_{\nu}(\boldsymbol{r},t) \psi_{\mu}^{*}(\boldsymbol{r},t) d\boldsymbol{r} dt = 0$$

- If $\,\nu\neq\mu$, the two modes will be orthogonal regardless of their spatial and temporal shapes
- When $\nu = \mu$, we will drop the polarization subscript
- When we say two "temporal modes" $s_1(t)$ and $s_2(t)$ are orthogonal, we will implicitly assume the mode functions being referred to have the same spatial modes and polarization

$$E(t) = a_1\phi_1(t) + a_2\phi_2(t) + 0\phi_3(t) + \dots \quad \square \qquad \text{Mode sorter} \qquad \square \qquad a_1\phi_1(t)$$
$$a_2\phi_2(t)$$



Maximum number of orthogonal modes

- Consider temporal modes, $\phi_k(t), k = 1, \ldots K$
- ...and their Fourier transforms,
- $\Phi_k(f) = \int \phi_k(t) e^{-2\pi j f t} dt$ $\phi_k(t) = \int \Phi_k(f) e^{2\pi j f t} df$ • How many (K) orthogonal modes $\phi_k(t)$ can be "fit into" a time-bandwidth product of T x W? i.e.,

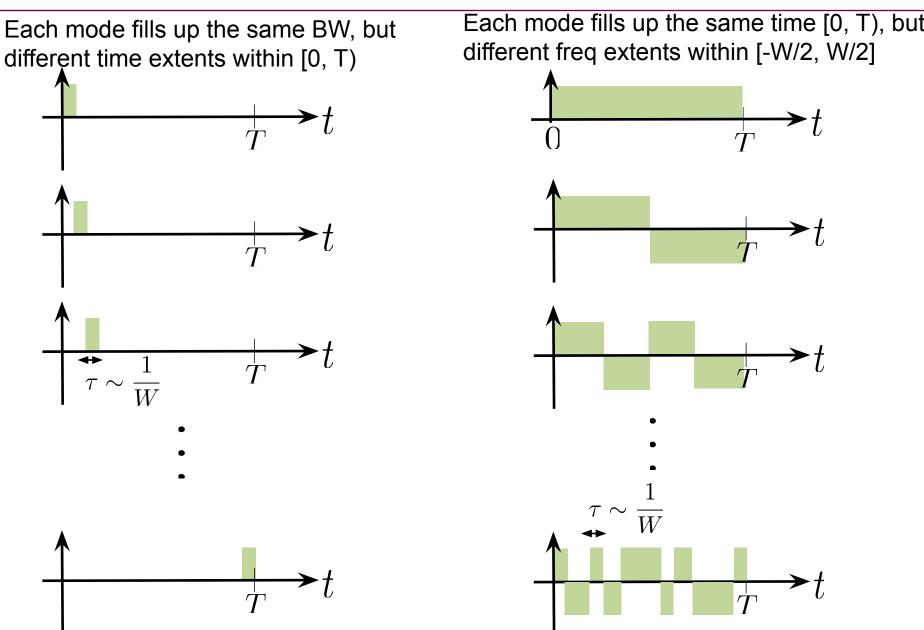
-
$$\phi_k(t)=0, t
otin[0,T)$$
 , and $\Phi_k(f)=0, f-f_0>$

– While ensuring orthogonality: $\int_{0}^{T} \phi_k(t)\phi_l^*(t)dt = \delta_{kl}$

- Answer: $K \approx WT$, and these optimal mode functions are "Prolate Spheroidal" functions
- All of above holds for spatial modes as well

Slepian, D. Prolate spheroidal wave functions, Fourier analysis and uncertainty — IV: Extension to many dimensions; Generalized prolate spheroidal functions Bell Syst. Tech. J., vol. 43, pp. 3009-3057, Nov. 1964.

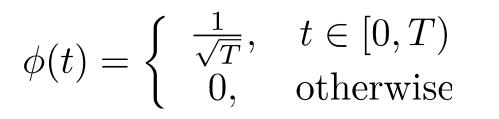
Some intuition: choices of WT modes

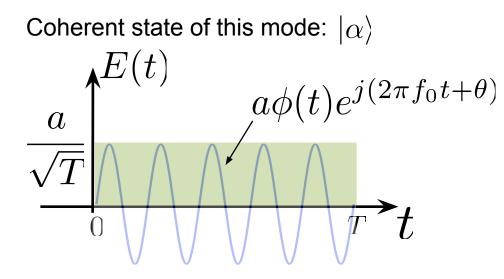


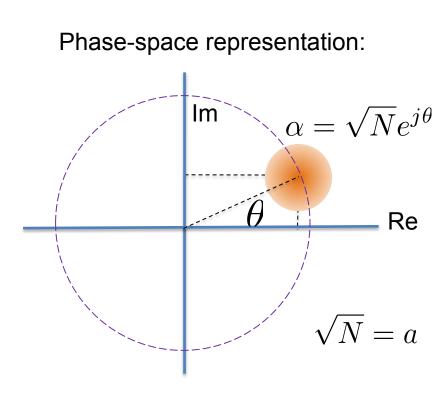
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Example: flat-top temporal mode



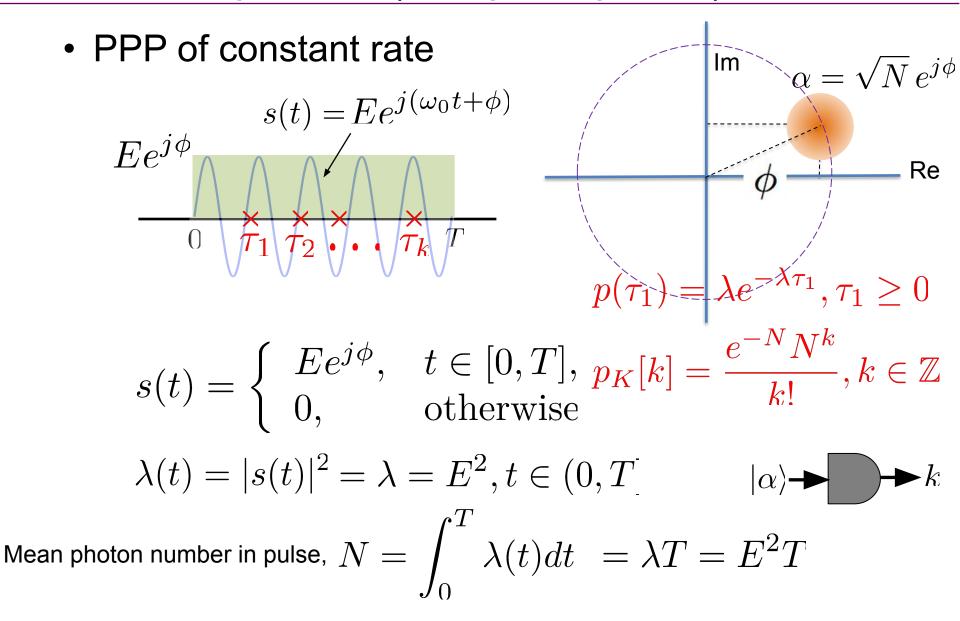




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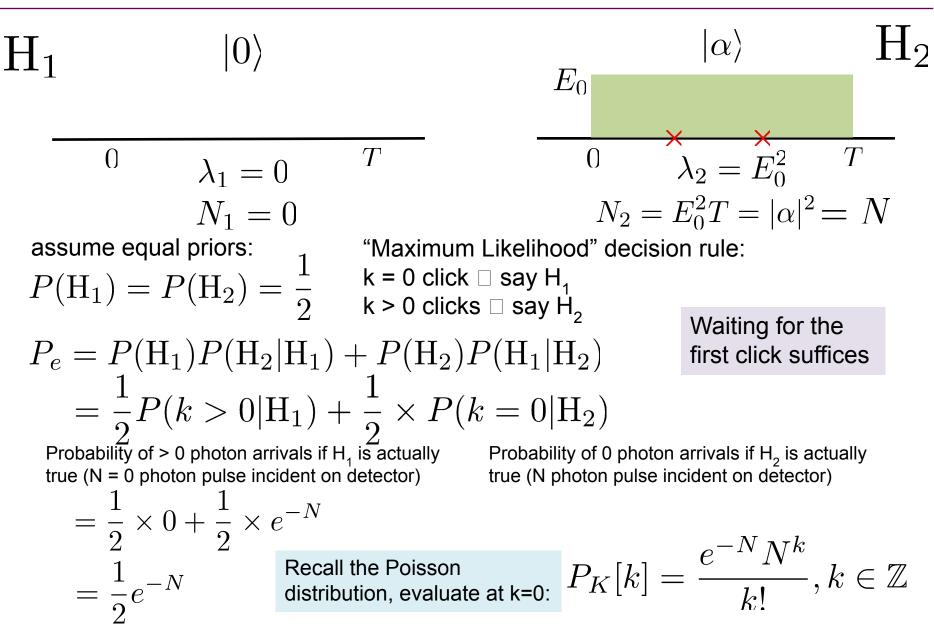
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Photon detection on a coherent state of Center for Quantum Networks (a "square pulse")



On-off keying (OOK) modulation

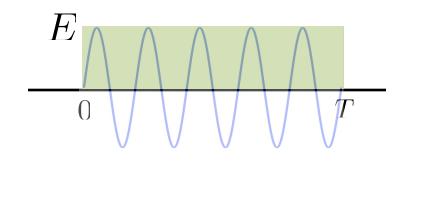


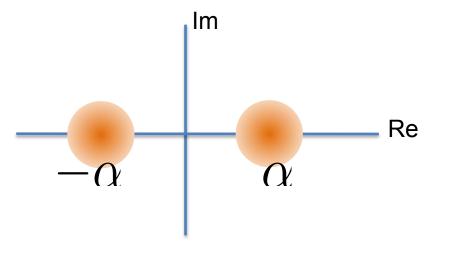


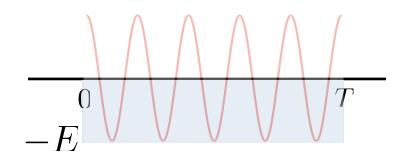


Binary phase shift keying (BPSK)

$$Ee^{j\phi}$$
 with, $\phi=\{0,\pi\}$







Mean photon number in the pulse is the same for either state:

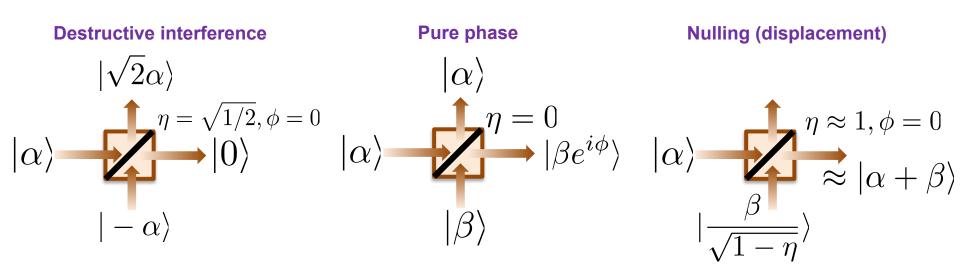
$$N = \left|\alpha\right|^2 = E^2 T$$



Interference (passive linear optics)

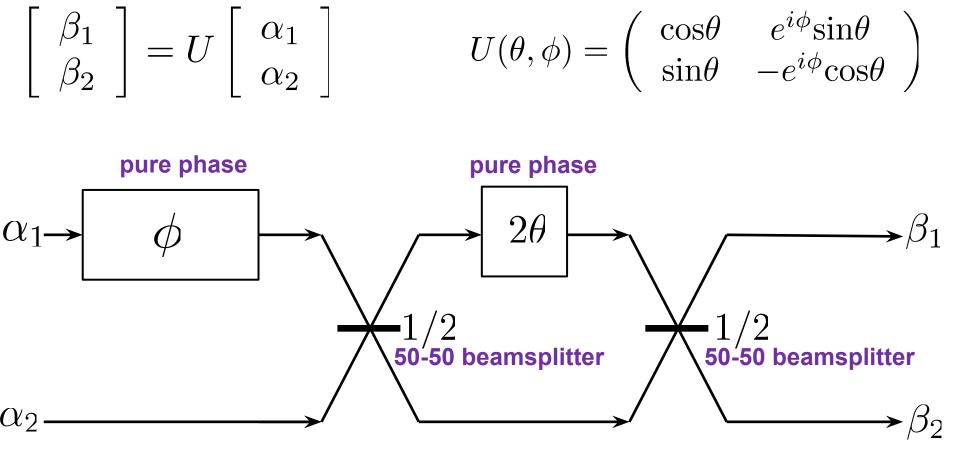
- $|\alpha_2\rangle$
 - Examples

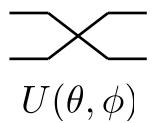
• Beamsplitter $\begin{vmatrix} \beta_1 \\ \beta_2 \end{vmatrix} = U \begin{bmatrix} \alpha_1 \\ \alpha_2 \end{bmatrix} \text{ Transmissivity, } \eta = \cos^2 \theta$ $|\beta_1\rangle \quad \eta \in (0,1)$ $\eta \in (0,1)$ $U(\theta,\phi) = \begin{pmatrix} \cos\theta & e^{i\phi}\sin\theta \\ \sin\theta & -e^{i\phi}\cos\theta \end{pmatrix}$ $\beta_1 = \sqrt{\eta}\alpha_1 + e^{i\phi}\sqrt{1-\eta}\alpha_2 \\ \beta_2 = \sqrt{1-\eta}\alpha_1 - e^{i\phi}\sqrt{\eta}\alpha_2 \quad \beta = U\alpha$





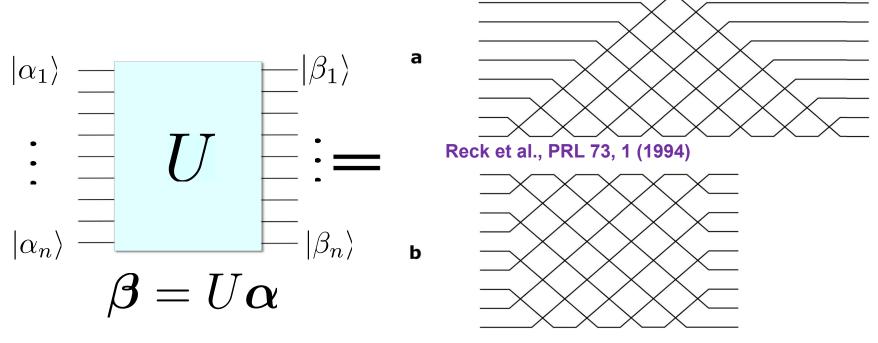
Mach-Zehnder interferometer (MZI)



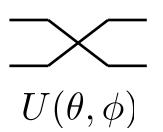


Arbitrary N-mode linear optical unitary

 Any N-by-N unitary U can be realized with M = N(N-1)/2 MZIs. Therefore, one needs tuning N(N-1) phases to realize any U

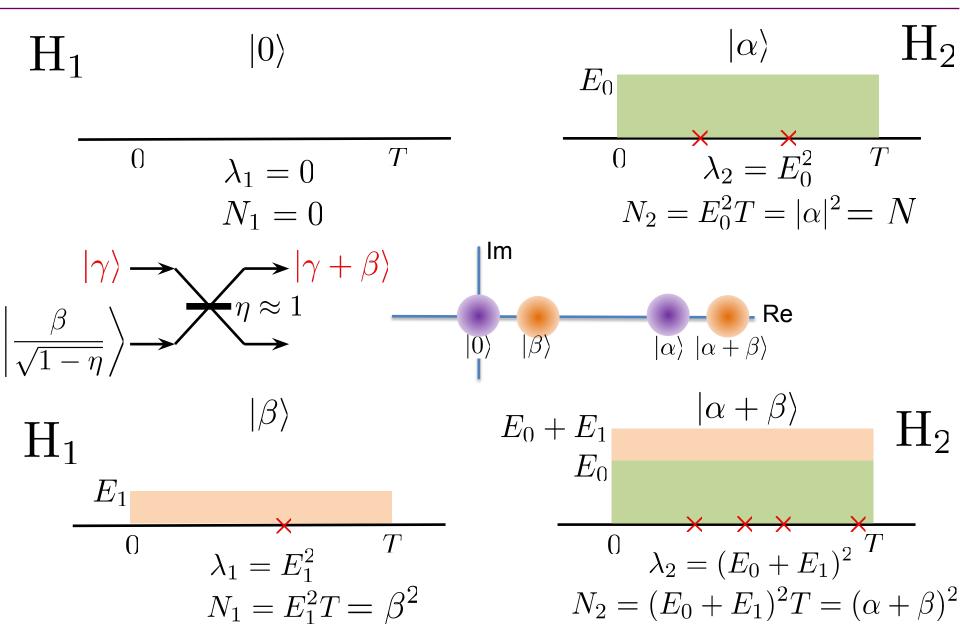


Clements et al., Optica 3 (12), 1460-1465 (2016)



On-off keying (OOK), Kennedy receiver NSF-ERC

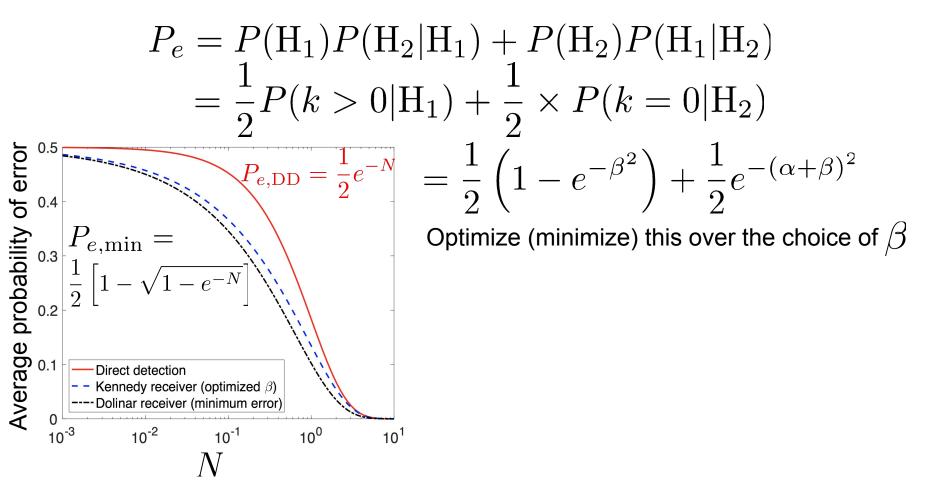
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P(error) for the Kennedy receiver

- Let us use the same decision rule
 - click = "on", no-click = "off"



Discriminating BPSK pulses $\{|-\alpha\rangle, \alpha\rangle$

• Kennedy (requires perfect amplitude and phase reference)

- Exact nulling
$$\{|-\alpha\rangle, |\alpha\rangle\} \longrightarrow \{|0\rangle, |2\alpha\rangle\}, |\alpha|^2 = N$$

R. Kennedy, MIT Research Laboratory for Electronics, Quarterly Progress Report 110, 219 – 225 (1972)

- Optimized nulling
$$\{|-\alpha\rangle, |\alpha\rangle\} \longrightarrow \{|\beta\rangle, |2\alpha + \beta\rangle\}$$

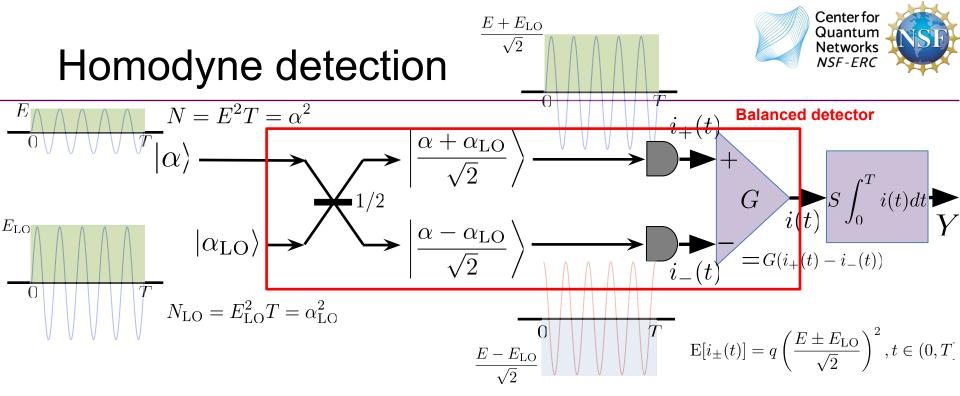
Takeoka and Sasaki, Phys. Rev. A 78, 022320 (2008)

$$P_{e}(N) = \min_{\beta} \left[\frac{1}{2} e^{-(2\alpha + \beta)^{2}} + \frac{1}{2} \left(1 - e^{-\beta^{2}} \right) \right] \qquad \beta \qquad 2\alpha + \beta$$

$$= \frac{1}{2} e^{-4N}, \ \beta = 0 \ \ \mbox{(exact nulling, suboptimal)}$$

Can we build a receiver for BPSK that does NOT require an amplitude reference?

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Assume for now that both $lpha, lpha_{
m LO}$ are real, and $N_{
m LO} \gg N$

$$Y = qGS(K_{+} - K_{-})$$

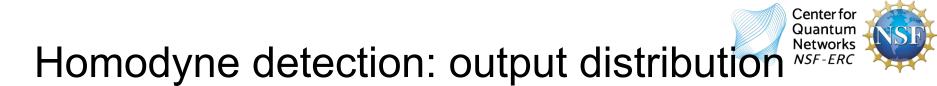
$$K_{+} \sim \text{Poisson}(N_{+}) \sim \mathcal{N}(N_{+}, N_{+}); N_{+} = \left|\frac{\alpha + \alpha_{\text{LO}}}{\sqrt{2}}\right|^{2}$$

$$K_{-} \sim \text{Poisson}(N_{-}) \sim \mathcal{N}(N_{-}, N_{-}); N_{-} = \left|\frac{\alpha - \alpha_{\text{LO}}}{\sqrt{2}}\right|^{2}$$

$$Y \sim \mathcal{N}(\mu, \sigma^2) \qquad \text{Poisson}(\lambda) \approx \text{Gaussian}(\lambda, \lambda), \lambda \gg 1$$

$$\mu = qGS(N_+ - N_-) = 2qGS\alpha\alpha_{\text{LO}} \qquad e^{-\lambda}\frac{\lambda^r}{r!} \simeq \frac{1}{\sqrt{2\pi\lambda}}e^{-\frac{(r-\lambda)^2}{2\lambda}}$$

$$\sigma^2 = (qGS)^2(N_+ + N_-) = (qGS)^2(\alpha^2 + \alpha_{\text{LO}}^2)$$



$$Y \sim \mathcal{N}(\mu, \sigma^2)$$

$$\mu = qGS(N_+ - N_-) = 2qGS\alpha\alpha_{\rm LO}$$

$$\sigma^2 = (qGS)^2(N_+ + N_-) = (qGS)^2(\alpha^2 + \alpha_{\rm LO}^2)$$

Pick the scaling constant S such that the mean $\mu = \alpha$

Problem 2: What is the distribution of Y?

A: $Y \sim \mathcal{N}(\alpha, 1/4)$ B: $Y \sim \mathcal{N}(\alpha, 1/2)$ C: $Y \sim \mathcal{N}(\alpha, 1)$ D: I do not know.



Break [5 minutes]

Homodyne detection (continued)

 $\alpha = \sqrt{N} e^{j\theta}$

Im

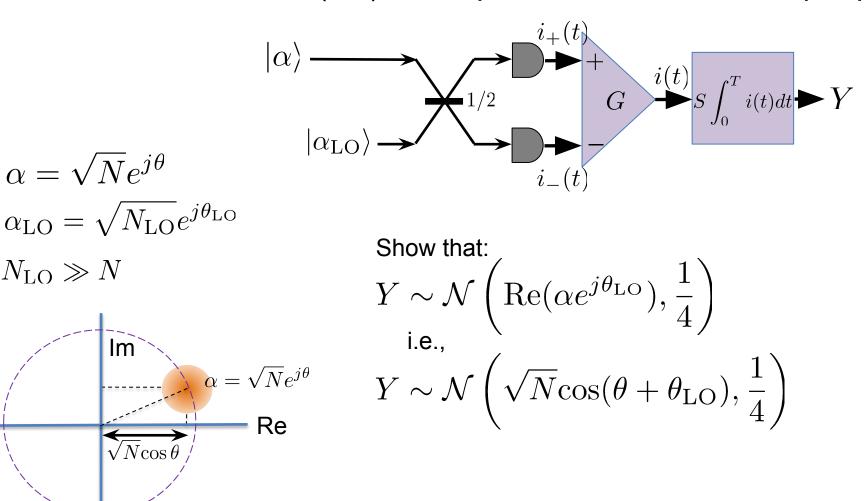
 $\sqrt{N}\cos\hat{\theta}$

 $N_{\rm LO} \gg N$

Local Oscillator (LO) has a phase offset with the input pulse

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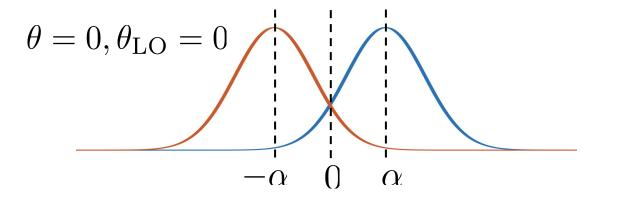
Homodyne detection



Black box description:

$$|\alpha\rangle \blacktriangleright \theta_{\rm LO} \rightharpoonup Y \sim \mathcal{N}\left(\operatorname{Re}(\alpha e^{j\theta_{\rm LO}}), \frac{1}{4}\right)$$
$$\alpha = \sqrt{N}e^{j\theta}$$

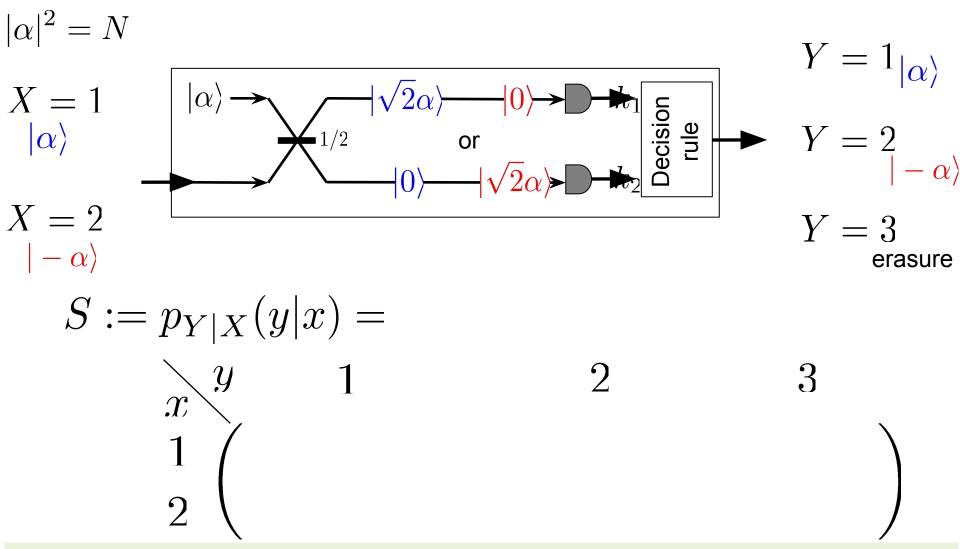
 $Y \sim \mathcal{N}\left(\mu, \sigma^2\right) \Rightarrow p_Y(y) = \frac{1}{\sqrt{2\pi}\,\sigma} e^{-(y-\mu)^2/2\sigma^2} \quad \text{Gaussian probability distribution}$



$$P_e = \frac{1}{2} \operatorname{erfc}(\sqrt{2N})$$

Unambiguous BPSK state discrimination

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If forced to make a decision, i.e. by mapping the erasure to one possible input, then the average probability of error (assume equal priors), $P_{e,\rm USD}=e^{-2N}$

Dual homodyne detection (heterodyne)

Recall black box description:

$$|\alpha\rangle \blacktriangleright \theta_{\rm LO} \blacktriangleright Y \sim \mathcal{N}\left(\operatorname{Re}(\alpha e^{j\theta_{\rm LO}}), \frac{1}{4}\right)$$

$$\sqrt{N}e^{j\theta}$$

 $-\alpha$

 $-j\alpha$

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Q-ary PSK, Q=4

 α

Re

$$\alpha = \sqrt{N}e^{j\theta} \xrightarrow{1}{2} \theta_{\text{LO}} \rightarrow Y_1 \sim \mathcal{N}\left(\operatorname{Re}\left(\frac{\alpha}{\sqrt{2}}e^{j\theta_{\text{LO}}}\right), \frac{1}{4}\right)$$

$$\alpha = \sqrt{N}e^{j\theta} \xrightarrow{1}{2} \theta_{\text{LO}} \xrightarrow{\pi}{2} \rightarrow Y_2 \sim \mathcal{N}\left(\operatorname{Re}\left(\frac{\alpha}{\sqrt{2}}e^{j(\theta_{\text{LO}}-\frac{\pi}{2})}\right), \frac{1}{4}\right)$$

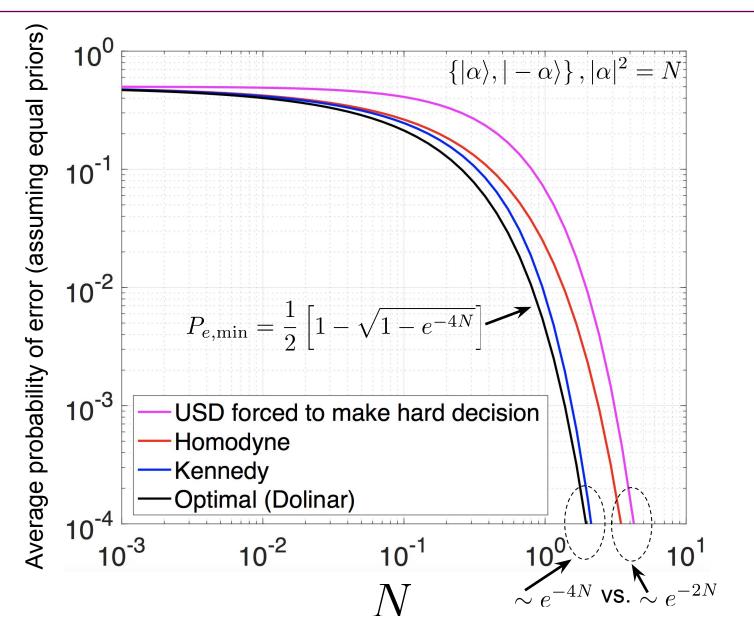
$$\frac{|\alpha|}{|\alpha|} \xrightarrow{j\alpha} \qquad j\alpha \qquad j\alpha \qquad \beta \in \mathcal{N}$$

Define:
$$X_1 = Y_1\sqrt{2}, X_2 = Y_2\sqrt{2}$$

 $\alpha =$

We can show: $X_1 \sim \mathcal{N}\left(\operatorname{Re}(\alpha e^{j\theta_{\mathrm{LO}}}), \frac{1}{2}\right)$ $X_2 \sim \mathcal{N}\left(\operatorname{Im}(\alpha e^{j\theta_{\mathrm{LO}}}), \frac{1}{2}\right)$

BPSK discrimination performance







Quantum states and measurements

- A (pure) state is described by a unit-norm column vector $|\psi\rangle$
- Von Neumann (projective) measurement on a state is described by a set of unit-norm orthonormal vectors $\{|w_k\rangle\}, \langle w_k|w_j\rangle = \delta_{kj}$
- If the state $|\psi\rangle$ is measured, the k-th outcome appears with probability, $p_k=|\langle w_k|\psi\rangle|^2$
- If two states $|\psi_1\rangle$ and $|\psi_2\rangle$ are orthogonal, i.e., $\langle \psi_1 | \psi_2 \rangle = 0$ measurement $\{|w_1\rangle = |\psi_1\rangle, |w_2\rangle = |\psi_2\rangle\}$ achieves $P_e = 0$



Mode $\phi(t)$, a quantum system, is excited in a coherent state $|lpha
angle, lpha\in\mathbb{C}$.

If we do ideal direct detection of mode $\,\phi(t),$ the total number of photons K is a Poisson random variable of mean N

Mode $\phi(t)$, a quantum system, is excited in a number state $|n
angle, n\in\{0,1,\dots,\infty\}$

If we do ideal direct detection of mode $\phi(t)$, the total number of photons K = n (exactly so; K is not a random variable).

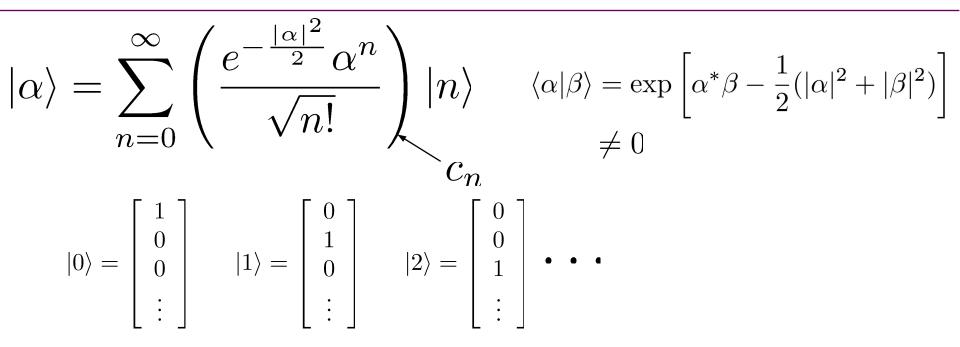
A mode of ideal laser light is in a coherent state. Number (Fock) state of a given mode is VERY hard to produce experimentally

There are infinitely many other types of "states" of the mode $\phi(t)$. Coherent state and Fock state are just two example class of states

 $|n
angle, n \in \{0, 1, \dots, \infty\}$ Fock states of a mode are special: they form an orthonormal basis that spans any general quantum state $|\psi
angle$ of that mode $\langle m|n
angle = \delta_{mn}$ and $|\psi
angle = \sum_{n=0}^{\infty} c_n |n
angle$, $\sum_{n=0}^{\infty} |c_n|^2 = 1$



Coherent state as a quantum state



Ideal photon detection is a von Neumann quantum measurement described by projectors, $\{|n\rangle\langle n|\}\,,n=0,1,\ldots,\infty$

Ideal direct detection on a coherent state $|\alpha\rangle$ produces outcome "n"

(i.e., n "clicks") with probability,
$$p_n = |\langle n | \alpha \rangle|^2 = |c_n|^2 = \frac{e^{-N}N^n}{n!}$$

Poisson detection statistics in a laser pulse is a result of the projection of the quantum state of the laser pulse—a coherent state—on to one of the Fock states

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Quantum description of light

- Complete description of an optical field is the quantum state of a set of mutually-orthogonal modes
- The most general (pure) state of a mode is an arbitrary superposition of number states of that mode, $|\psi\rangle = \sum_{n=0}^{\infty} c_n |n\rangle$ where $\langle m|n\rangle = \delta_{mn}$, and , $\sum_{n=0}^{\infty} |c_n|^2 = 1$
- Ideal direct detection: $\{|w_k^{n=0}\} \equiv \{|k\rangle\}$ -- number states
- Examples of pure state of a mode:

- Coherent state
$$|\alpha\rangle = \sum_{n=0}^{\infty} \left(\frac{e^{-\frac{|\alpha|^2}{2}}\alpha^n}{\sqrt{n!}}\right)|n\rangle \quad , \alpha \in \mathbb{C}$$

- Number (Fock) state, $|n\rangle, n \in \{1, 2, \ldots\}$
- Cat state: $|\psi_{\pm}\rangle = \mathcal{N}_{\pm} (|\alpha\rangle \pm |-\alpha\rangle)$ • $\mathcal{N}_{\pm} = 1/\sqrt{2(1 \pm e^{-2|\alpha|^2})}$ ensures $\langle \psi_{\pm} | \psi_{\pm} \rangle = 1$, and $\langle \psi_{-} | \psi_{-} \rangle = 1$ • $\langle \psi_{\pm} | \psi_{-} \rangle = 0$ (so, the two cat states can be used to encode a qubit)



Photon detection on cat states

$$|lpha
angle=e^{-rac{1}{2}|lpha|^2}\sum_{n=0}^\inftyrac{lpha^n}{\sqrt{n!}}|n
angle \qquad |-lpha
angle=e^{-rac{1}{2}|-lpha|^2}\sum_{n=0}^\inftyrac{(-lpha)^n}{\sqrt{n!}}|n
angle$$

 $\ket{\mathrm{cat}_e} \propto \ket{lpha} + \ket{-lpha}$

$$|\mathrm{cat}_e
angle \propto 2e^{-rac{1}{2}|lpha|^2} \left(rac{lpha^0}{\sqrt{0!}}|0
angle + rac{lpha^2}{\sqrt{2!}}|2
angle + rac{lpha^4}{\sqrt{4!}}|4
angle + \ldots
ight)$$

Not Poisson distribution

$$\ket{ ext{cat}_o} \propto \ket{lpha} - \ket{-lpha}$$

$$\left | \mathrm{cat}_o
ight
angle \propto 2 e^{-rac{1}{2} |lpha|^2} \left (rac{lpha^1}{\sqrt{1!}} |1
angle + rac{lpha^3}{\sqrt{3!}} |3
angle + rac{lpha^5}{\sqrt{5!}} |5
angle + \ldots
ight)$$



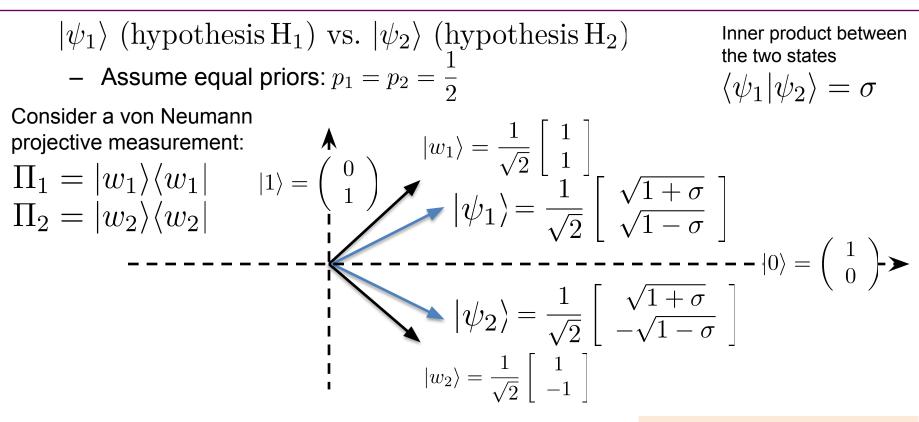
Entangled states

- Product state of two modes can be written as: $|\psi_1\rangle |\psi_2\rangle$ with, $|\psi_1\rangle = \sum_{n=0}^{\infty} a_n |n\rangle$ and $|\psi_2\rangle = \sum_{n=0}^{\infty} b_n |n\rangle$
 - Tensor product: Each state "lives in its own Hilbert space"
 - Example of product state of 2 modes:
 - Two coherent states, $|lpha_1
 angle|lpha_2
 angle$
 - Coherent state and a Fock state, $|\alpha\rangle|k\rangle$
 - Two cat states, $|\psi_{\pm}\rangle|\psi_{\pm}\rangle$, where $|\psi_{\pm}\rangle = \mathcal{N}_{\pm} (|\alpha\rangle \pm |-\alpha\rangle)$.
- Entangled state of two modes cannot be written

as $|\psi_1\rangle|\psi_2\rangle$ A two-mode entangled state is: $|\psi\rangle = \sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} c_{n_1,n_2}|n_1\rangle|n_2\rangle$, $\sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} |c_{n_1,n_2}|^2 = 1$ - Example (NOON state): $|\psi\rangle = \frac{|n\rangle|0\rangle + |0\rangle|n\rangle}{\sqrt{2}}$



Binary state discrimination



$$P_{e} = P(\mathbf{H}_{1})P(\mathbf{H}_{2}|\mathbf{H}_{1}) + P(\mathbf{H}_{2})P(\mathbf{H}_{1}|\mathbf{H}_{2})$$
$$= \frac{1}{2} |\langle w_{2}|\psi_{1}\rangle|^{2} + \frac{1}{2} |\langle w_{1}|\psi_{2}\rangle|^{2}$$

This measurement happens to be the one that minimizes the average error probability

Minimum probability of error for binary state discrimination



$$P_e = \frac{1}{2} |\langle w_2 | \psi_1 \rangle|^2 + \frac{1}{2} |\langle w_1 | \psi_2 \rangle|^2$$
$$|\psi_1 \rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} \sqrt{1+\sigma} \\ \sqrt{1-\sigma} \end{bmatrix} \qquad |w_1 \rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$
$$|\psi_2 \rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} \sqrt{1+\sigma} \\ -\sqrt{1-\sigma} \end{bmatrix} \qquad |w_2 \rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

Problem 3: What is the minimum probability of error in discriminating $|\psi_1\rangle$ and $|\psi_2\rangle$ given their inner product, $\langle \psi_1 | \psi_2 \rangle = \sigma$? (assume σ is real) A: $P_e = \sigma$ B: $P_e = [1 - \sqrt{1 - \sigma^2}]/2$ C: $P_e = [1 - \sigma]/2$ D: I do not know.



MPE decision among M pure states

referr

as

- M-ary ensemble $\{p_i, |\psi_i\rangle\}; i = 1, 2, ..., M$
 - Pairwise inner products (Gram matrix), $\sigma_{ij} = \langle \psi_i | \psi_j \rangle$
- Yuen-Kennedy-Lax (YKL) [1975] conditions for MPE
 - Projective measurement, $\{|w_i\rangle\}, j = 1, 2, \dots, M$
 - Relative inner products (aligning measurement vectors in the M-dimensional space spanned by the M pure states):

$$x_{ij} = \langle w_i | \psi_j \rangle$$

- YKL conditions for minimum average probability of error

(1)
$$p_m x_{km} x_{mm}^* = p_k x_{kk} x_{mk}^*$$

Sometimes
referred to
as
"Helstrom
bound"
(1) $p_m x_{km} x_{mm}^* = p_k x_{kk} x_{mk}^*$
(2) $\sum_{k=1}^{M} x_{kj} x_{ki}^* = \sigma_{ij}$
(3) $\sum_{i=1}^{M} x_{jj} |w_j\rangle \langle \psi_j | \ge p_i |\psi_i\rangle \langle \psi_i |, \forall i$
(3) $\sum_{i=1}^{M} x_{jj} |w_j\rangle \langle \psi_j | \ge p_i |\psi_i\rangle \langle \psi_i |, \forall i$
(5) $M(M+1)/2 \text{ non-linear simultaneous equations: solve for \mathcal{X}_{ij}
 $P_{e,\min} = 1 - \sum_{i=1}^{M} p_i |x_{ii}|^2$
(6) (Check for uniqueness of solution)$



BPSK – minimum probability of error

$$\alpha \rangle = \sum_{n=0}^{\infty} \left(\frac{e^{-\frac{|\alpha|^2}{2}} \alpha^n}{\sqrt{n!}} \right) |n\rangle$$

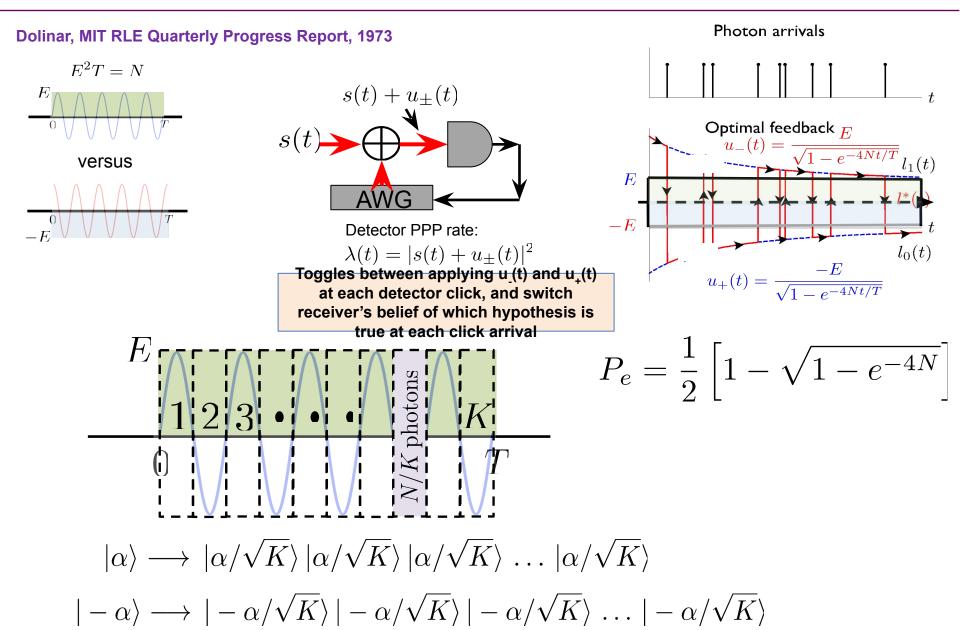
- The inner product of two coherent states $|\alpha\rangle$ and $|\beta\rangle$, $\langle \alpha |\beta \rangle = \exp \left[\alpha^* \beta - \frac{1}{2} (|\alpha|^2 + |\beta|^2) \right]$
- Discriminating BPSK coherent states, $\{|\alpha\rangle, |-\alpha\rangle\}, \alpha \in \mathbb{R}$

- Inner product, $\sigma = \langle \alpha | - \alpha \rangle = e^{-2N}, |\alpha|^2 = N$ Im $P_{e.\min} = \frac{1}{2} \left[1 - \sqrt{1 - |\sigma|^2} \right] = \frac{1}{2} \left[1 - \sqrt{1 - e^{-4N}} \right]$ $- \frac{1}{2} \left[1 - \sqrt{1 - e^{-4N}} \right]$ $- \frac{1}{2} \left[1 - \sqrt{1 - e^{-N/2}} \right]$ $- \frac{1}{2} \left[1 - \sqrt{1 - e^{-N/2}} \right]$ $- \frac{1}{2} \left[1 - \sqrt{1 - e^{-N/2}} \right]$

This calculation of minimum error probability using the quantum picture was easy. How do we design a receiver that will achieve this?

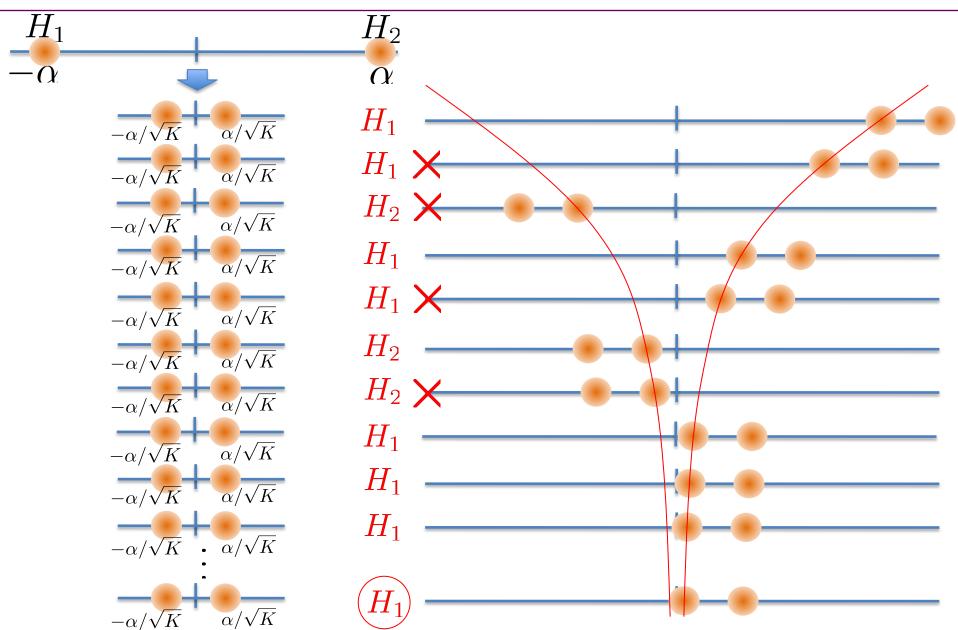
Dolinar's receiver: optimal binary detection

Center for

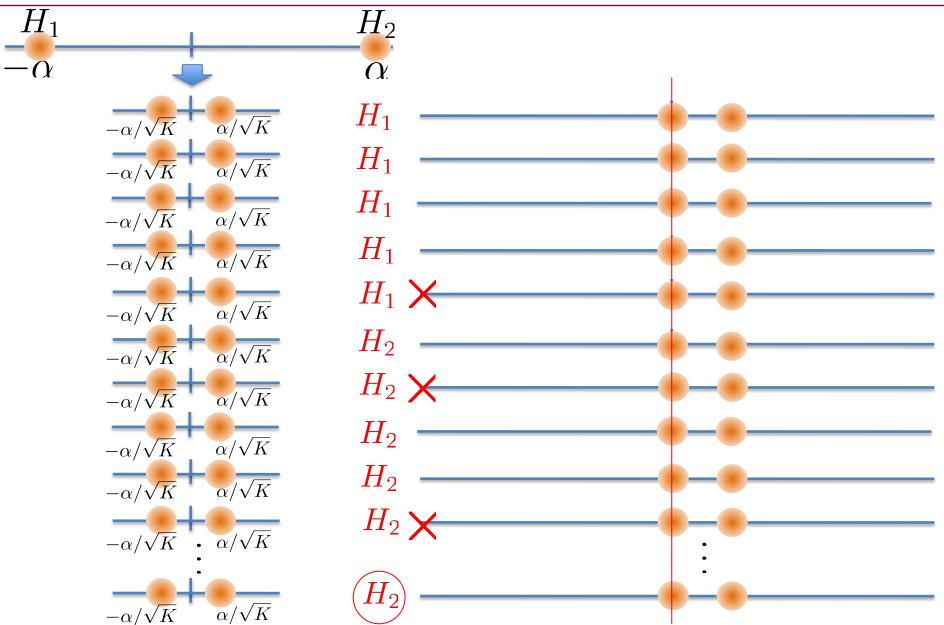


Slicing interpretation of Dolinar's receivers

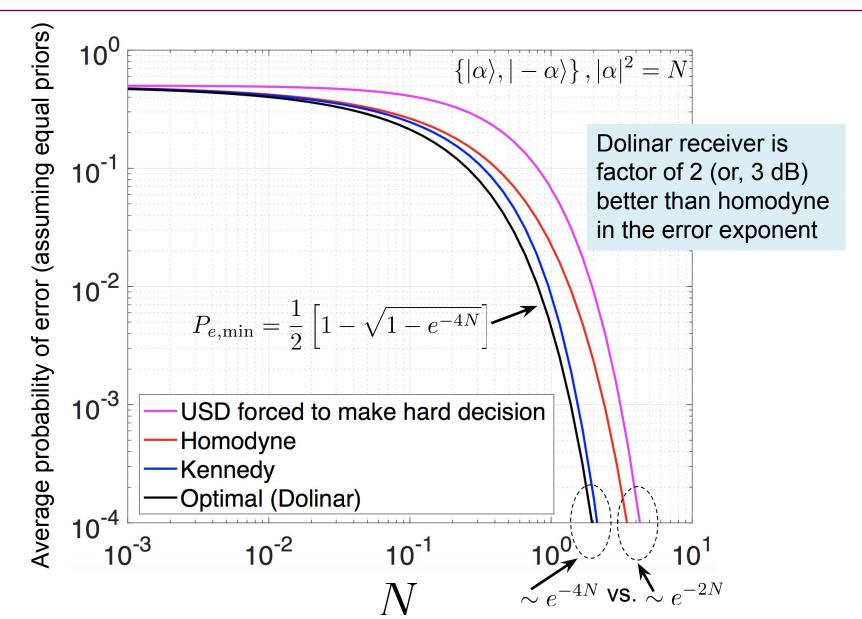
Center for Quantum



Slicing interpretation of Kennedy's receiver (exact nulling version)

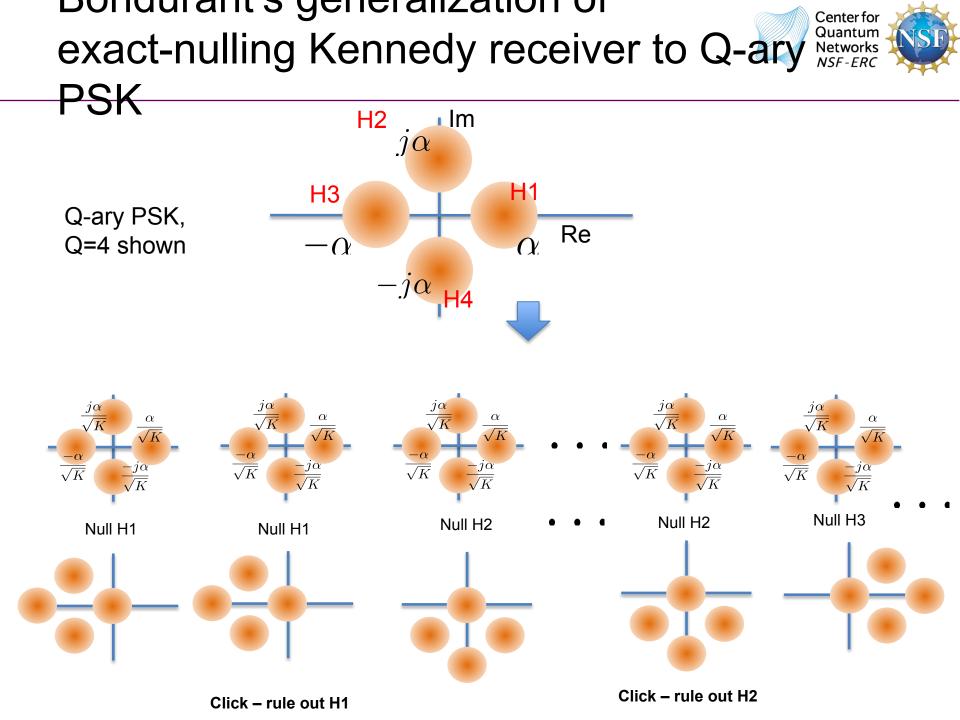


BPSK discrimination performance



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Optimized nulling Bondurant receiver: the equivalent of Dolinar for Q-ary PSK



C. R. Müller and C. Marquardt, New J. Phys. 17 032003 (2015)

They analyze performance of the receiver with all three detector imperfections we discussed

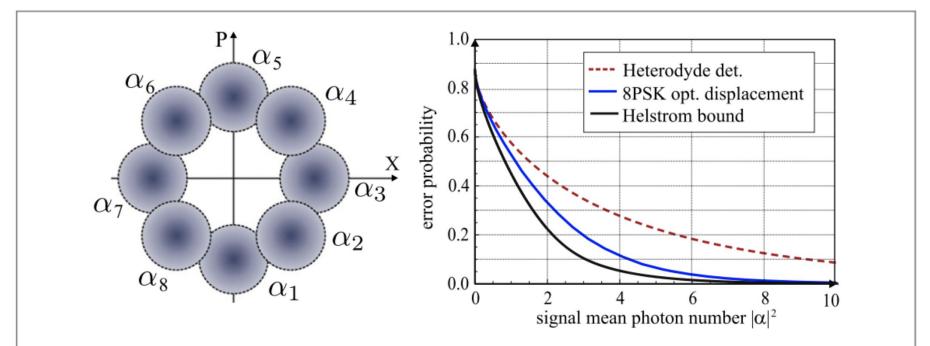


Figure 5. (Left) Illustration of the 8PSK alphabet in phase space. (Right) Error probabilities of the 8PSK optimized displacement receiver (with Bayesian probing) compared to the standard quantum limit (heterodyne detection) and the Helstrom bound. Just like for the QPSK alphabet, the standard quantum limit is outperformed for any signal power.

The "quantum" limit of minimum probability of error cannot be attained by this Dolinar-like receiver strategy. A structured receiver to attain this minimum error probability remains an open problem even for 3 given states of a laser pulse!



Ternary discrimination example

- Consider $|\psi_1\rangle = |0\rangle, \ |\psi_2\rangle = |\alpha\rangle, \ |\psi_3\rangle = |-\alpha\rangle$
 - with α real, equal priors, $N = |\alpha|^2$; what measurement would you use to distinguish these?
- Gram matrix, $((\Gamma))_{ij} = \gamma_{ij} = \langle \psi_i | \psi_j \rangle$

$$\Gamma = \left(\begin{array}{ccc} 1 & x & x \\ x & 1 & x^4 \\ x & x^4 & 1 \end{array}\right), x = e^{-N/2}$$
State geometry
Measurement geometry

• Measurement matrix, $((X))_{ij} = x_{ij} = \langle w_i | \psi_j \rangle$ $X = \begin{pmatrix} a & d & d \\ b & c & e \\ b & e & c \end{pmatrix}$ $|\psi_1\rangle = |0\rangle$ $|\psi_1\rangle = |0\rangle$ $|\psi_2\rangle = |\alpha\rangle\rangle$

Minimum probability of error



YKL conditions

$$\sum_{k=1}^{M} x_{kj} x_{ki}^{*} = \gamma_{ij} \begin{cases} a^{2} + 2b^{2} = 1 \\ d^{2} + c^{2} + e^{2} = 1 \\ ad + b(c + e) = x \\ d^{2} + 2ce = x^{4} \end{cases}$$

$$p_m x_{km} x_{mm}^* = p_k x_{kk} x_{mk}^* \left\{ ab = cd \right\}$$

- Solution:

$$a = \left[2xd + (1 - x^2)\sqrt{1 + x^4 - 2d^2} \right] / (1 + x^4)$$

$$b = \left[x\sqrt{1 + x^4 - 2d^2} - d(1 - x^2) \right] / (1 + x^4)$$

$$c = \frac{1}{2} \left[\sqrt{1 + x^4 - 2d^2} + \sqrt{1 - x^4} \right]$$

$$e = \frac{1}{2} \left[\sqrt{1 + x^4 - 2d^2} - \sqrt{1 - x^4} \right]$$

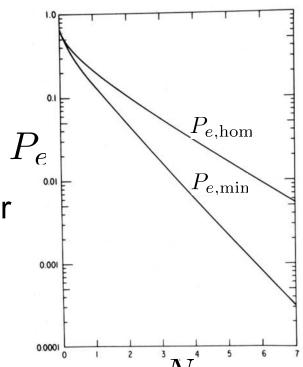
• Substitute these into f(d) = ab - cd, and solve using Newton's method for the value of d, s.t. f(d) = 0. Evaluate average min probability of error, $P_{e,\min} = 1 - \frac{1}{3}(a^2 + 2c^2)$ and plot as fn. of N



Comparison of MPE with Homodyne

- Evaluate the error probability attained by an ideal homodyne detection receiver: $X \sim \left(x_i, \frac{1}{4}\right), x_i = 0, -\alpha, \alpha$ - If $|X| < -\alpha/2$, pick $|-\alpha\rangle$, else if $|X| > \alpha/2$, pick $|\alpha\rangle$, else pick $|0\rangle$
- Show that, $P_{e,\text{hom}} = \frac{4}{3} \operatorname{erfc}(\sqrt{N})$. Let us plot this. Asymptotic limit (N large)
- - For N large, $P_{e,\mathrm{hom}}\sim e^{-N/2}$, whereas $P_{e,\min} \sim e^{-N}$
- Kennedy like receiver (sequential-nulling) Kennedy receiver outperforms homodyne, and attains the optimal exponent, exp(-N)

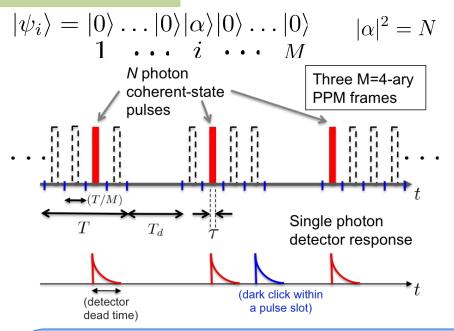
Plot courtesy: Quantum Detection and Estimation Theory, Carl Helstrom, 1976



Another M-ary example: pulse position modulation (PPM)



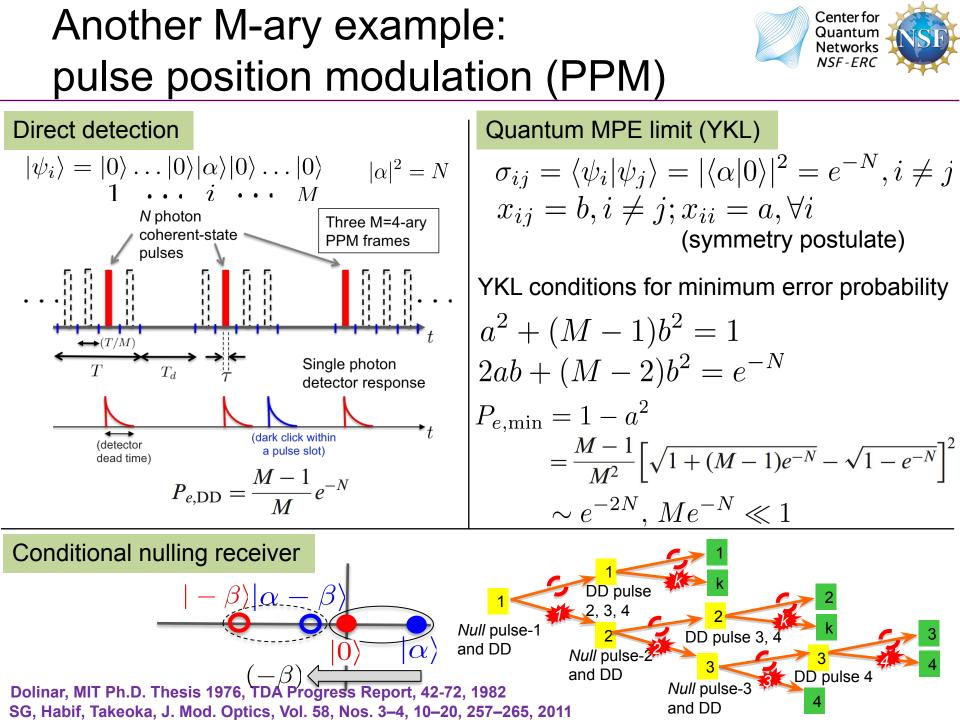
Direct detection



Problem 4: What is the probability of error in discriminating the M-ary PPM codewords achieved by ideal direct (photon) detection on each pulse slot?

A:
$$P_e = e^{-N}$$

B: $P_e = e^{-2N}/M$
C: $P_e = [(M - 1)/M]e^{-N}$
D: I do not know.



PPM demodulation using the Conditional Pulse Nulling (CPN) receiver

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$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + Me^{-N} - 1 \right]$$

$$\sim e^{-2N}, Me^{-N} \ll 1$$

$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + Me^{-N} - 1 \right]$$

$$\sim e^{-2N}, Me^{-N} \ll 1$$

$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + Me^{-N} - 1 \right]$$

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$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + Me^{-N} - 1 \right]$$

$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + 1 \right]$$

$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + 1 \right]$$

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$$P_{e,\text{CPN}} = \frac{1}{M} \left[(1 - e^{-N})^M + 1 \right]$$

$$P_{e,\text{CPN}} =$$

Chen, Habif, Dutton, Lazarus and SG, Nature Photonics 6, 374-379 (2012)

Receiver design that exactly attains the quantum limit is not known.

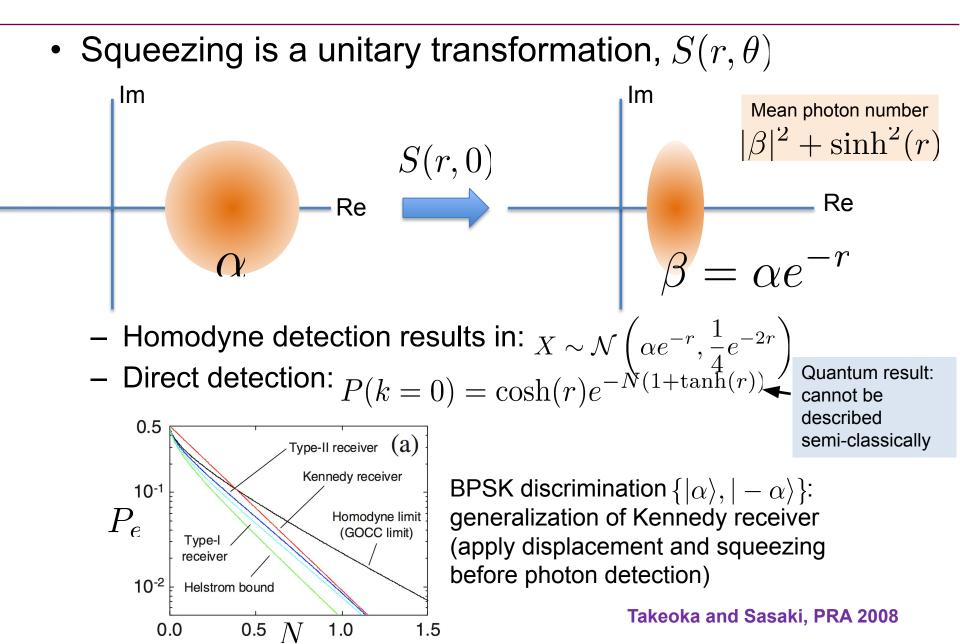
Universal quantum processing

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- We know how to calculate the minimum probability of error for discriminating any M coherent states. Yet, we don't know optimal structured receiver designs for M > 2
- The M = 2 case (Dolinar receiver) was special
- So far, we have been playing with linear optics (circuits of beamsplitters and phases) and direct (photon) detection. These are NOT universal resources
- Adding "squeezing" to our toolbox will make it *universal*
- We will learn about squeezing in Module 2

Squeezing







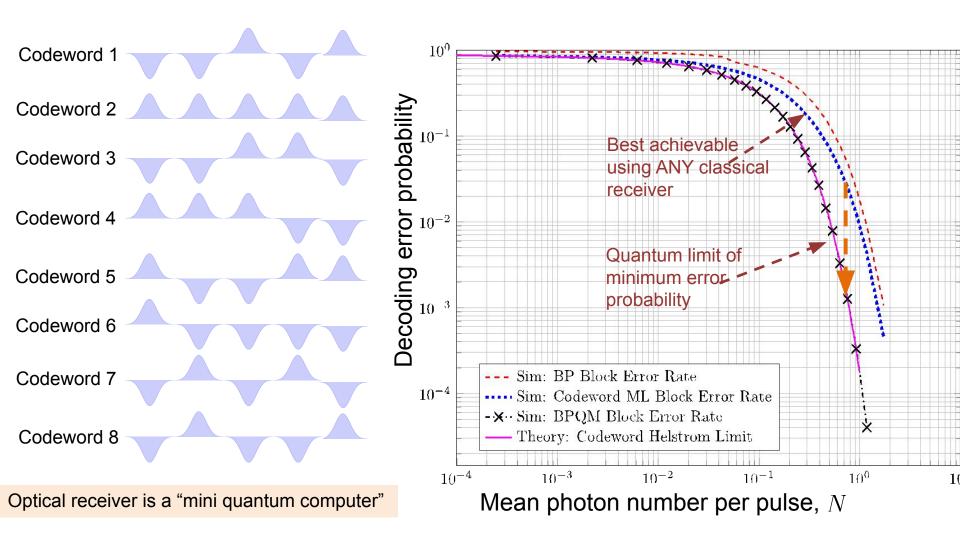
General design of an optimal receiver

- Is there a receiver strategy that uses adaptive application of squeezing (not just displacement) on small slices of the coherent state pulses, and photon detection attain arbitrary M-ary MPE state discrimination? [Open problem]
- Instead of an all-optical design, what if we can map each of the BPSK coherent states to a qubit first, i.e., $|\alpha\rangle \longrightarrow |\psi_1\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} \sqrt{1+\sigma} \\ \sqrt{1-\sigma} \\ \sqrt{1-\sigma} \end{bmatrix}$ $|-\alpha\rangle \longrightarrow |\psi_2\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} \sqrt{1+\sigma} \\ \sqrt{1+\sigma} \\ -\sqrt{1-\sigma} \end{bmatrix}$

$$\sigma = \langle -\alpha | \alpha \rangle = e^{-2N}, \, N = |\alpha|^2$$

– Then use *quantum computing* on those qubits?

Quantum limited receiver design



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Silva, SG, Dutton, Phys. Rev. A 87, 052320 (2013) Rengaswamy, Seshadreesan, SG, Pfister, Nature npj Quantum Inf 7, 97 (2021) Delaney, Seshadreesan, MacCormack, Galda, SG, Narang, Phys. Rev. A 106, 032613 (2022)

Demonstration of quantum advantage



PHYSICAL REVIEW A 106, 032613 (2022)

Demonstration of a quantum advantage by a joint detection receiver for optical communication using quantum belief propagation on a trapped-ion device

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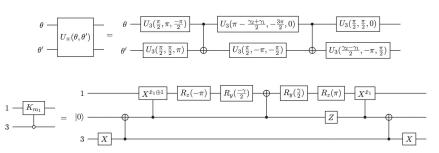
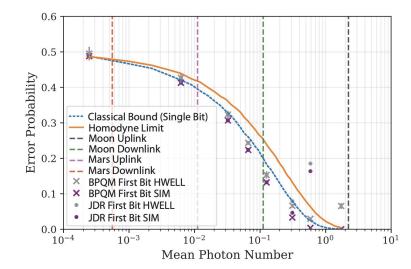


FIG. 5. Both decompositions for BPQM full decoder components. (a) U gate decomposition, where U_3 is the QISKIT rotation gate and γ_1 and γ_2 are defined in Eq. (A9). (b) K_m gate decomposition, utilizing ancilla qubit 3.



Communication capacity

(a

ancilla mode

code)

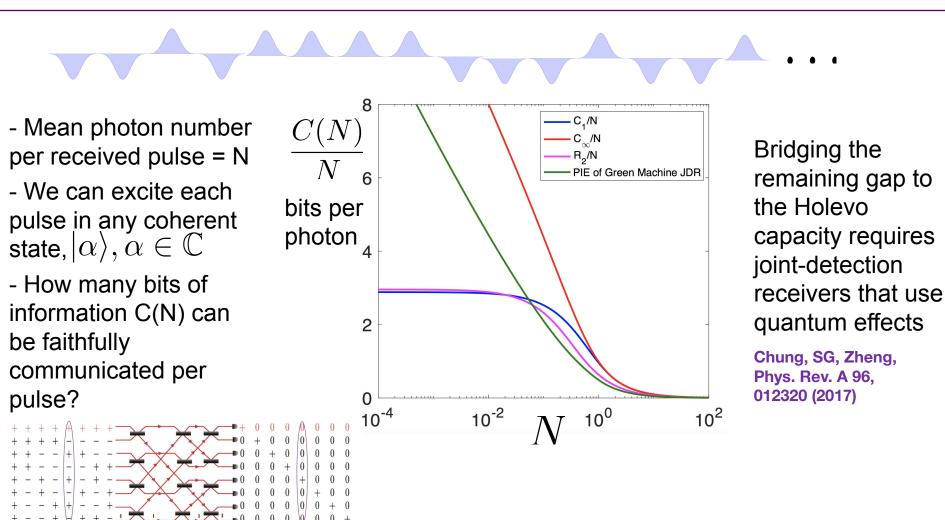
(BPSK Hadamard

Unitary operation

(the Walsh transform)

(Pulse Position Modulation)





SG, Phys. Rev. Lett. 106, 240502 (2011)

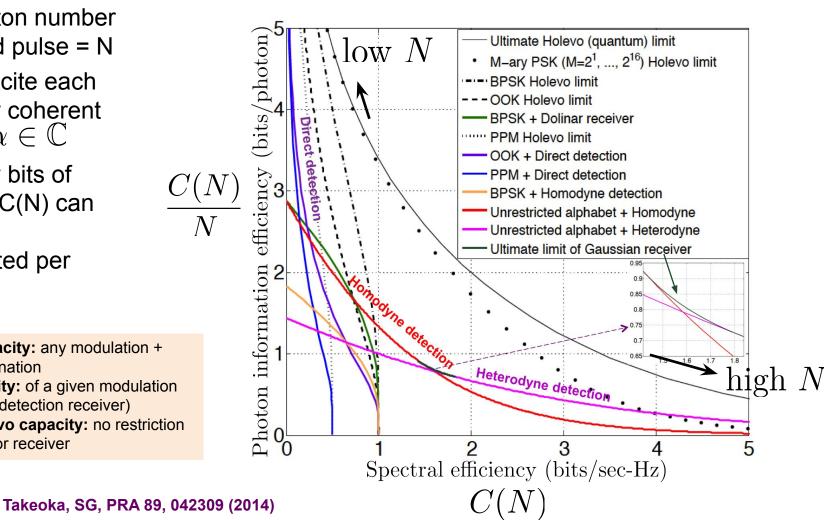
Communication capacity





- Mean photon number per received pulse = N
- We can excite each pulse in any coherent state, $|\alpha\rangle, \alpha \in \mathbb{C}$
- How many bits of information C(N) can be faithfully communicated per pulse?

Shannon capacity: any modulation + receiver combination Holevo capacity: of a given modulation (optimum joint detection receiver) Ultimate Holevo capacity: no restriction of modulation or receiver



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Module 1: concluding remarks

- Laser light pulses undergo "wave-like" interference (through beamsplitters), much like ripples of water in a pond
- Laser light field cannot be precisely measured
- Detecting photons on a laser light pulse produces a random number of "clicks" with a Poisson distribution
- Quantum representation of a laser light pulse is a "coherent state"
 --- this representation helps us quantify "best" receivers (that minimize probability of error, for example), even without knowing how to build such as optimal receiver
- Just using semiclassical tools (interference in a beamsplitter based circuit, and Poisson-noise-limited photon detection), one cannot attain the quantum limit of receiver performance
- Optimal receiver designs require "quantum" processing of laser light – either all-optically using non-classical transformations of light (e.g., using squeezing) or first coupling the laser-light pulses into qubits, followed by processing them in a quantum computer



Break [5 minutes]



Outline:

- 1. Introduction to very basic quantum mechanics.
- 2. Beam-splitters: Classical v Quantum Inputs.
- 3. Gaussian Multiport interferometers.
- 4. General Gaussian Transformations and application to Gaussian Boson Sampling.

Quantum Systems



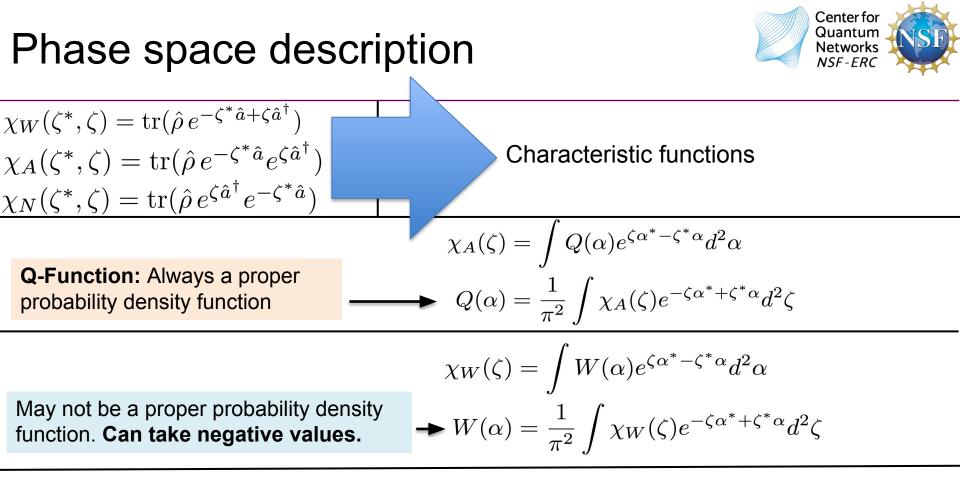
Essentially, Quantum optics is applied linear algebra.



CV System: We encode information in states of light.

Transformation of states Networks NSF-ERC Phase Squeezing Displacement $\hat{U}_{eta}{=}\exp\!\left(eta\hat{a}^{\dagger}{-}eta^{*}\hat{a} ight)$

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Always a proper probability density function *when it exists*. The states for which a proper P function exists are classical states.

$$\chi_N(\zeta) = \int P(\alpha) e^{\zeta \alpha^* - \zeta^* \alpha} d^2 \alpha$$
$$\bullet P(\alpha) = \frac{1}{\pi^2} \int \chi_N(\zeta) e^{-\zeta \alpha^* + \zeta^* \alpha} d^2 \zeta$$



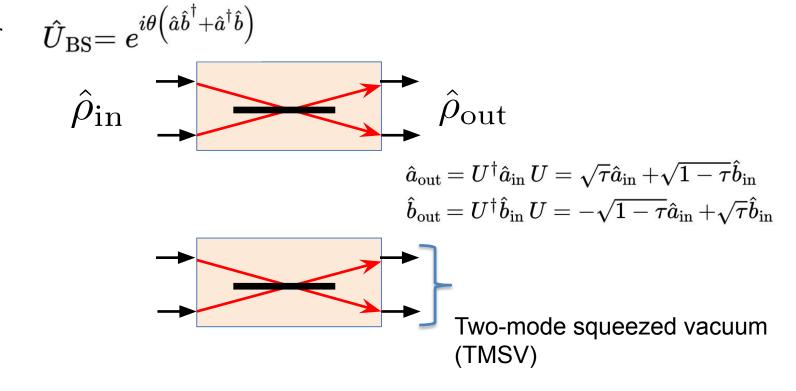


$$\begin{split} \chi_W(\zeta^*,\zeta) &= \operatorname{tr}(\hat{\rho} \, e^{-\zeta^* \hat{a} + \zeta \hat{a}^\dagger}) & \chi_W^U(\zeta^*,\zeta) = \operatorname{tr}(\hat{U} \hat{\rho} \hat{U}^\dagger e^{-\zeta^* \hat{a} + \zeta \hat{a}^\dagger}) & \chi_W^U(\zeta^*,\zeta) = \operatorname{tr}(\hat{\rho} \, e^{-\zeta^* \hat{U}^\dagger \hat{a} \hat{U} + \zeta \hat{U}^\dagger \hat{a}^\dagger \hat{U}}) \\ \chi_A(\zeta^*,\zeta) &= \operatorname{tr}(\hat{\rho} \, e^{-\zeta^* \hat{a}} e^{\zeta \hat{a}^\dagger}) & \swarrow \chi_A^U(\zeta^*,\zeta) = \operatorname{tr}(\hat{U} \hat{\rho} \hat{U}^\dagger e^{-\zeta^* \hat{a}} e^{\zeta \hat{a}^\dagger}) & \chi_A^U(\zeta^*,\zeta) = \operatorname{tr}(\hat{\rho} \, e^{-\zeta^* \hat{U}^\dagger \hat{a} \hat{U}} e^{\zeta \hat{U}^\dagger \hat{a}^\dagger \hat{U}}) \\ \chi_N(\zeta^*,\zeta) &= \operatorname{tr}(\hat{\rho} \, e^{\zeta \hat{a}^\dagger} e^{-\zeta^* \hat{a}}) & \chi_N^U(\zeta^*,\zeta) = \operatorname{tr}(\hat{U} \hat{\rho} \hat{U}^\dagger e^{\zeta \hat{a}^\dagger} e^{-\zeta^* \hat{a}}) & \chi_N^U(\zeta^*,\zeta) = \operatorname{tr}(\hat{\rho} \, e^{\zeta \hat{U}^\dagger \hat{a}^\dagger \hat{U}^\dagger} e^{-\zeta^* \hat{U}^\dagger \hat{a} \hat{U}}) \end{split}$$

Gaussian Unitary Operator: Any unitary operator that maps Gaussian states to Gaussian states.

Two-mode transformation: Beam splitter

Beam splitter

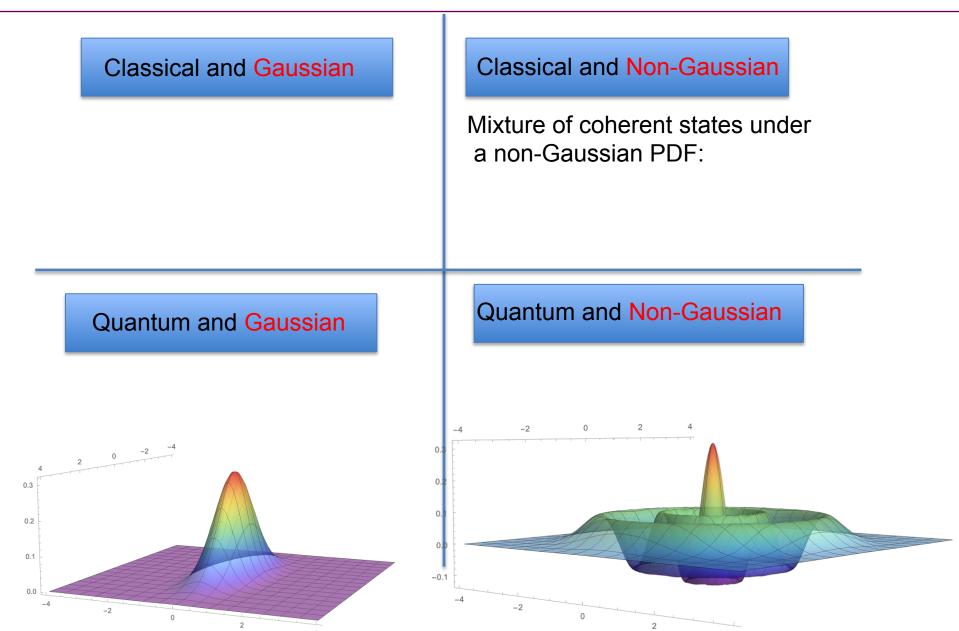


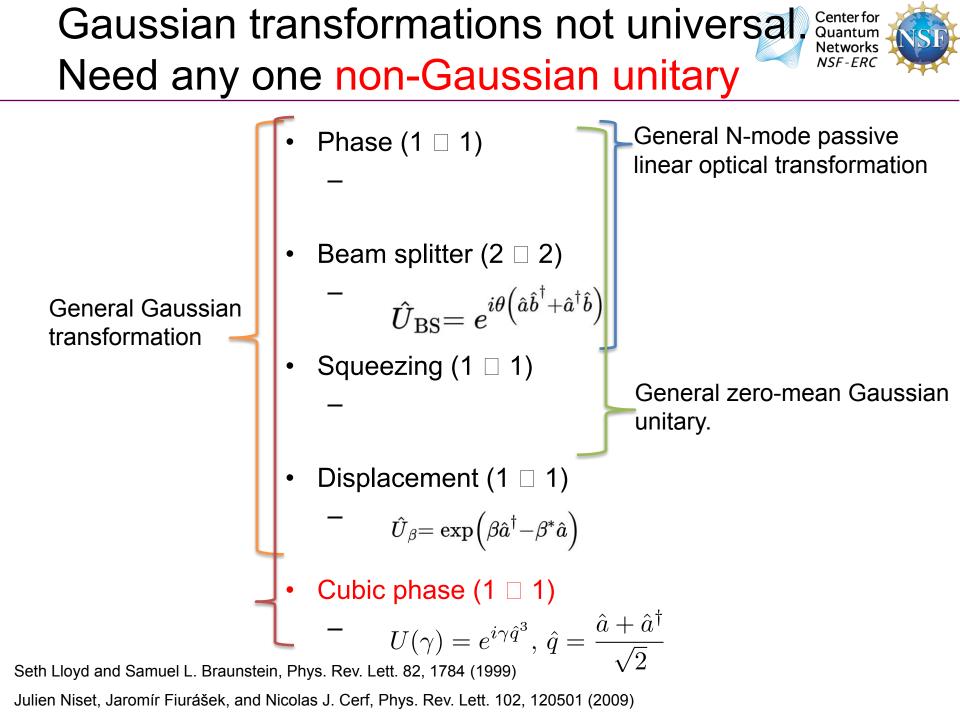
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Networks

Examples of States Light



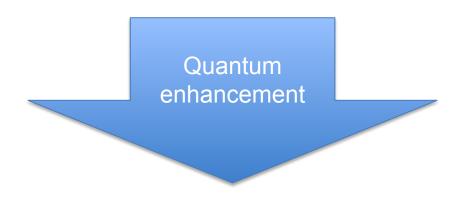






Purely Quantum properties

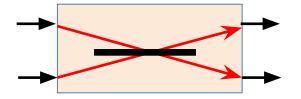
- 1. Squeezing (Gaussian Quantum resources).
- 2. Entanglement.
- 3. Non-Gaussian states and measurements.



- Sensing
- Communications.
- Universal Quantum Computation

Example on beam splitter calculations



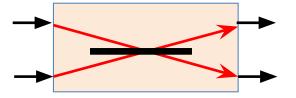


 $egin{aligned} \hat{a}_{ ext{out}} &= U^{\dagger} \hat{a}_{ ext{in}} \, U = \sqrt{ au} \hat{a}_{ ext{in}} + \sqrt{1 - au} \hat{b}_{ ext{in}} \ \hat{b}_{ ext{out}} &= U^{\dagger} \hat{b}_{ ext{in}} \, U = -\sqrt{1 - au} \hat{a}_{ ext{in}} + \sqrt{ au} \hat{b}_{ ext{in}} \end{aligned}$

 $\label{eq:linear}$

Example on beam splitter calculations (cont.)





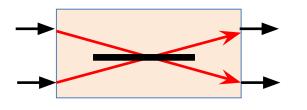
 $egin{aligned} \hat{a}_{ ext{out}} &= U^\dagger \hat{a}_{ ext{in}} \, U = \sqrt{ au} \hat{a}_{ ext{in}} + \sqrt{1- au} \hat{b}_{ ext{in}} \ \hat{b}_{ ext{out}} &= U^\dagger \hat{b}_{ ext{in}} \, U = -\sqrt{1- au} \hat{a}_{ ext{in}} + \sqrt{ au} \hat{b}_{ ext{in}} \end{aligned}$



A coherent state input in a beam splitter **cannot** produce entanglement: The resulting state is a product of two coherent states (while the total mean photon number is conserved).

More examples on beam splitter calculations

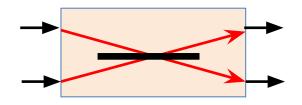




 $egin{aligned} \hat{a}_{ ext{out}} &= U^{\dagger} \hat{a}_{ ext{in}} \, U = \sqrt{ au} \hat{a}_{ ext{in}} + \sqrt{1 - au} \hat{b}_{ ext{in}} \ \hat{b}_{ ext{out}} &= U^{\dagger} \hat{b}_{ ext{in}} \, U = -\sqrt{1 - au} \hat{a}_{ ext{in}} + \sqrt{ au} \hat{b}_{ ext{in}} \end{aligned}$



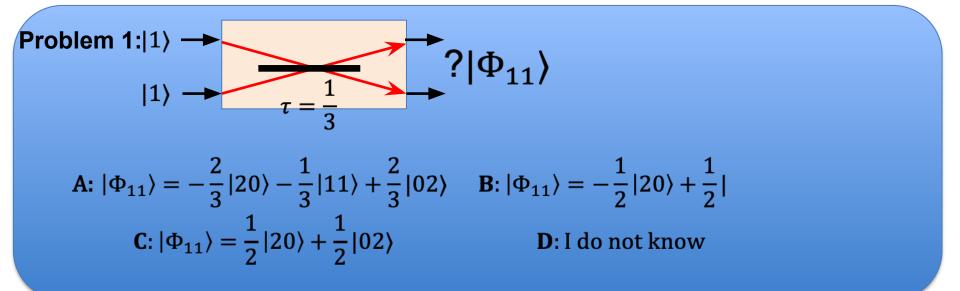






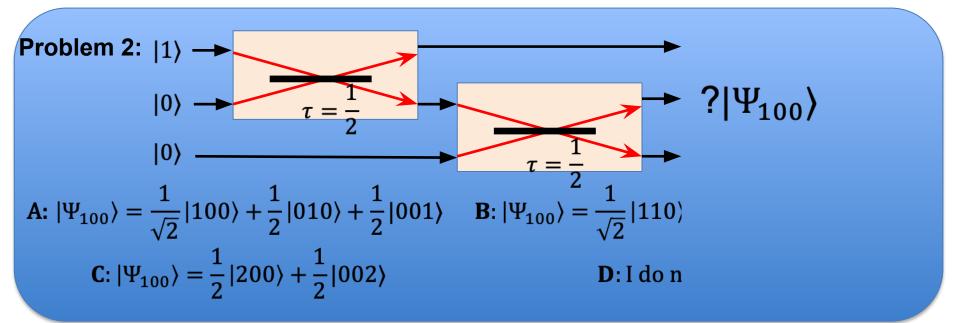
Problems (7 minutes)





Problem







Break [5 minutes]

Introducing multi-mode Gaussian interferometers



We have seen that the beam-splitter transforms the annihilation and creation as:

$$egin{aligned} \hat{a}_{ ext{out}} &= U^{\dagger} \hat{a}_{ ext{in}} \, U = \sqrt{ au} \hat{a}_{ ext{in}} + \sqrt{1 - au} \hat{b}_{ ext{in}} \ \hat{b}_{ ext{out}} &= U^{\dagger} \hat{b}_{ ext{in}} \, U = -\sqrt{1 - au} \hat{a}_{ ext{in}} + \sqrt{ au} \hat{b}_{ ext{in}} \end{aligned}$$

$$\hat{a}_{ ext{out}}^{\dagger} = \sqrt{ au} \hat{a}_{ ext{in}}^{\dagger} + \sqrt{1- au} \hat{b}_{ ext{in}}^{\dagger} \ \hat{b}_{ ext{out}}^{\dagger} = -\sqrt{1- au} \hat{a}_{ ext{in}}^{\dagger} + \sqrt{ au} \hat{b}_{ ext{in}}^{\dagger}$$

$$egin{aligned} \hat{b}_1 &= \sqrt{ au} \hat{a}_1 \!+\! \sqrt{1- au} \hat{a}_2 \ \hat{b}_2 &= -\sqrt{1- au} \hat{a}_1 \!+\! \sqrt{ au} \hat{a}_2 \end{aligned}$$

$$egin{aligned} \hat{b}_{1}^{\dagger} &= \sqrt{ au} \hat{a}_{1}^{\dagger} + \sqrt{1 - au} \hat{a}_{2}^{\dagger} \ \hat{b}_{2} &= -\sqrt{1 - au} \hat{a}_{1}^{\dagger} + \sqrt{ au} \hat{a}_{2}^{\dagger} \end{aligned}$$



Multimode Gaussian Unitary evolution

For example: The beam splitter corresponds to:
$$V = \begin{pmatrix} \sqrt{\tau} & \sqrt{1-\tau} & 0 & 0 \\ 0 & 0 & \sqrt{\tau} & \sqrt{1-\tau} \\ 0 & 0 & -\sqrt{1-\tau} & \sqrt{\tau} \end{pmatrix}$$



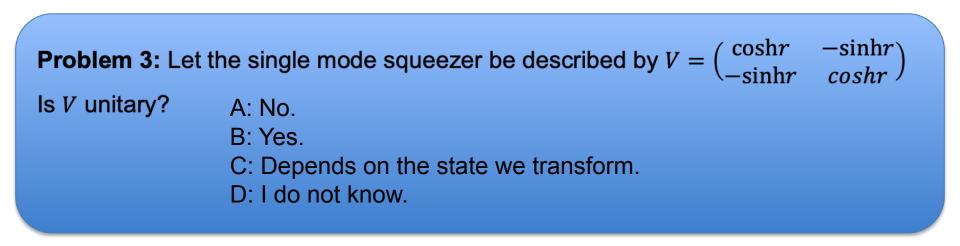


Multimode Gaussian Unitary evolution

$$egin{pmatrix} \hat{b}\ \hat{b}^\dagger \end{pmatrix} = Vigin{pmatrix} \hat{a}\ \hat{a}^\dagger \end{pmatrix}, \ ig(\hat{b}^\dagger \ \hat{b}ig)igg(\hat{b}\ \hat{b}^\dagger \end{pmatrix} = ig(\hat{a}^\dagger \ \hat{a}ig)V^\dagger Vig(\hat{a}\ \hat{a}^\dagger \end{pmatrix} \Rightarrow 2\hat{b}^\dagger \hat{b} + 1 = ig(\hat{a}^\dagger \ \hat{a}ig)V^\dagger Vig(\hat{a}\ \hat{a}^\dagger \end{pmatrix}$$







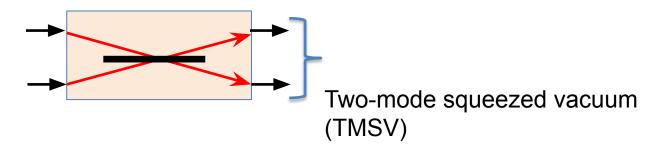
Problem



Problem 4: Let the single mode squeezed vacuum state $\widehat{U}_{squeezer}|0\rangle = |0;r\rangle = \frac{1}{\sqrt{\cosh n}} \sum_{n=0}^{\infty} \frac{\sqrt{(2n)!}}{n!} \left(-\frac{1}{2} \tanh r\right)^n |2n\rangle$ What is the mean photon number of said state? Reminder: Mean photon number $\overline{n} = \langle 0; r | \widehat{n} | 0; r \rangle$ and $\widehat{n} | n \rangle = n | n \rangle$ A: $\overline{n} = 0$. B: $\overline{n} = \sinh^2 r$. C: $\overline{n} = \frac{1}{2} \sinh^2 r$. D: I do not know.

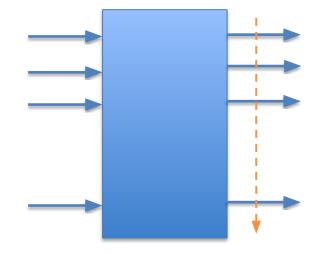


With the math you've learned you can calculate:



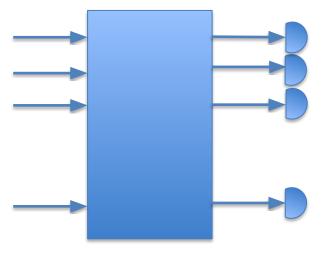
The TMSV is the maximally entangled two-mode Gaussian state.

A more general transformation:



The probability of any PNR pattern





$$P_{ec{n}} = rac{\left|\mathcal{I}_{n_{1}\dots n_{N}}
ight|^{2}}{\prod_{i=1}^{N}n_{i}!2^{n_{i}}\cosh r_{i}}$$
Andrew J. Pizzimenti, Joseph M. Lukens, Hsuan-Hao Lu, Nicholas A. Peters, Saikat Guha, and Christos N. Gagatsos, Phys. Rev. A **104**, 062437 (2021)
 $\mathcal{I}_{ec{n}} = \int d^{N}q_{lpha}d^{N}p_{lpha}R(ec{x}_{lpha})\prod_{i=1}^{N}(q_{lpha_{i}}+ip_{lpha_{i}})^{n_{i}} = \left\langle f_{1}^{n_{1}}\dots f_{N}^{n_{N}}
ight
angle, f_{i} = q_{lpha_{i}}+ip_{lpha_{i}}$
 $R(ec{x}_{lpha}) = rac{1}{(2\pi)^{N}}e^{-rac{1}{2}ec{x}_{lpha}^{T}\mathcal{H}ec{x}_{lpha}}$
 \mathcal{H}

The probability of any PNR pattern: Hafnians



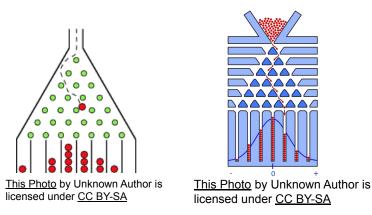
The difficult part in calculating he probability of a photon-number pattern, is the term:

$$\mathcal{I}_{ec{n}} = \int d^N q_lpha d^N p_lpha R(ec{x}_lpha) \prod_{i=1}^N (q_{lpha_i} + i p_{lpha_i})^{n_i} = ig\langle f_1^{n_1} \dots f_N^{n_N}ig
angle, f_i = q_{lpha_i} + i p_{lpha_i}$$

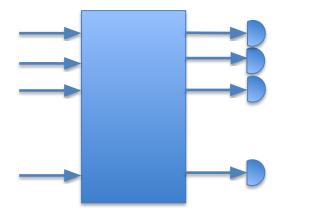
Hafnian of matrix F: $\mathcal{I}_{\vec{n}} = \begin{cases} 0 & \Sigma = \text{odd}, \\ \text{Hf}(F) & \Sigma = \text{even} \end{cases} \quad \Sigma = \sum_{i=1}^{N} n_i$

Sampling from a distribution



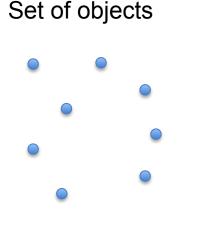


In general, a (classical) computer program that samples from a given distribution would scale exponentially with the size of the requested pattern (e.g. photon-number pattern). Therefore, we program the Hafnian into a quantum-optical circuit, i.e., a *quantum Galton board* (a.k.a. Gaussian Boson Sampler) that performs the sampling job.



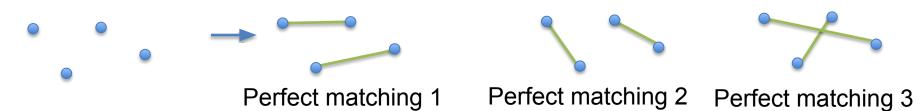
Hafnians and perfect matchings





Perfect matching: Connect the objects with a line so that any object is used only once (no more than one line can start/end from any point).

Example:



Hafnians and perfect matchings

How does the previous example relate to the calculation of the Hafnian?





Isserlis' theorem



$$P_{ec{n}} = rac{ig|\mathcal{I}_{n_1\dots n_N}ig|^2}{\prod_{i=1}^N n_i! 2^{n_i} {
m cosh}\, r_i} \hspace{1cm} \mathcal{I}_{ec{n}} = \int d^N q_lpha d^N p_lpha R(ec{x}_lpha) \prod_{i=1}^N (q_{lpha_i} + i p_{lpha_i})^{n_i} = ig\langle f_1^{n_1}\dots f_N^{n_N}ig
angle, f_i = q_{lpha_i} + i p_{lpha_i}$$

$$\langle g_1g_2\cdots g_\Sigma
angle = \sum_{p\in P_\Sigma^2}\prod_{\{i,j\}\in p}\langle g_ig_j
angle$$

Module 2: Concluding remarks

Center for Quantum Networks NSF-ERC

- Quantum states of light can:
 - Possess counter-intuitive properties
 - Give quantum advantage

- Necessary tools to manipulate all the different protocols:
 - Linear algebra.
 - Building intuition.

Quantum supremacy (advantage)



- Google, US Nature volume 574, pages 505–510 (2019)
 - IQP: a random circuit based sampler [53-qubit circuit of depth 20, with 430 two-qubit and 1,113 single-qubit gates. Classical simulation estimate ~ 10,000 years based on a Schrödinger-Feynman simulation that trades off space for time, whereas Sycamore processor takes about 200 seconds
- Xanadu, Canada Nature volume 606, pages 75–81 (2022)
 - Gaussian Boson Sampling: 216 squeezed modes entangled with three-dimensional connectivity; it would take ~ 9,000 years for the best available algorithms and supercomputers to produce, using exact methods, a single sample from the programmed distribution, whereas Borealis requires only 36 µs
- USTC, China Science, Vol 370, Issue 6523 pp. 1460-1463, 3 Dec 2020
 - Gaussian Boson Sampling: 50 squeezed modes in a 100-mode interferometer; measured a sampling rate that is about 10¹⁴-fold faster than using state-of-the-art classical simulation strategies and supercomputers (2.5 billion years as opposed to 200 s)

Center for Quantum Networks NSF-ERC

What are random samplers good for?

- Quantum generative adversarial networks (arXiv:1804.08641)
- Neural networks and other machine learning algorithms powered by continuous-variable quantum states and processes (arXiv:1806.06871, arXiv:1912.08278, arXiv:2001.03622)
- Graph isomorphism and graph similarity testing (arXiv:1810.10644)
- Molecular vibronic spectra for quantum chemistry calculations (arXiv:1912.07634)
- Quantum adiabatic optimization algorithm or QAOA (arXiv:1902.00409)
- Producing samples from hard-to-generate stochastic point processes (arXiv:1906.11972)
- Multi-variate optimization (arXiv:1909.02108)

Random samplers (Boson sampling, IQP, GBS, etc.) are not universal quantum computers. But they are believed to be strictly more powerful than classical computers



- The building blocks we covered in Modules 1 and 2 (linear optics, coherent states, multi-photon states, squeezed states, photon detection, homodyne detection) in principle suffice to design "quantum-optimal" transmitters, processors, computers, receivers, for all applications in photonics-based information processing
- Yet, we don't know how to "put together" these building blocks to attain best performance in most applications of photonic information processing!



Center for Quantum Networks

Course Evaluation Survey

We value your feedback on all aspects of this short course. Please go to the link provided in the Zoom Chat or in the email you will soon receive to give your opinions of what worked and what could be improved.

CQN Winter School on Quantum Networks

Funded by National Science Foundation Grant #1941583













